SAT-Based Verification with IC3: Foundations and Demands

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Induction

Foundation of verification for 40+ years (Floyd, Hoare)

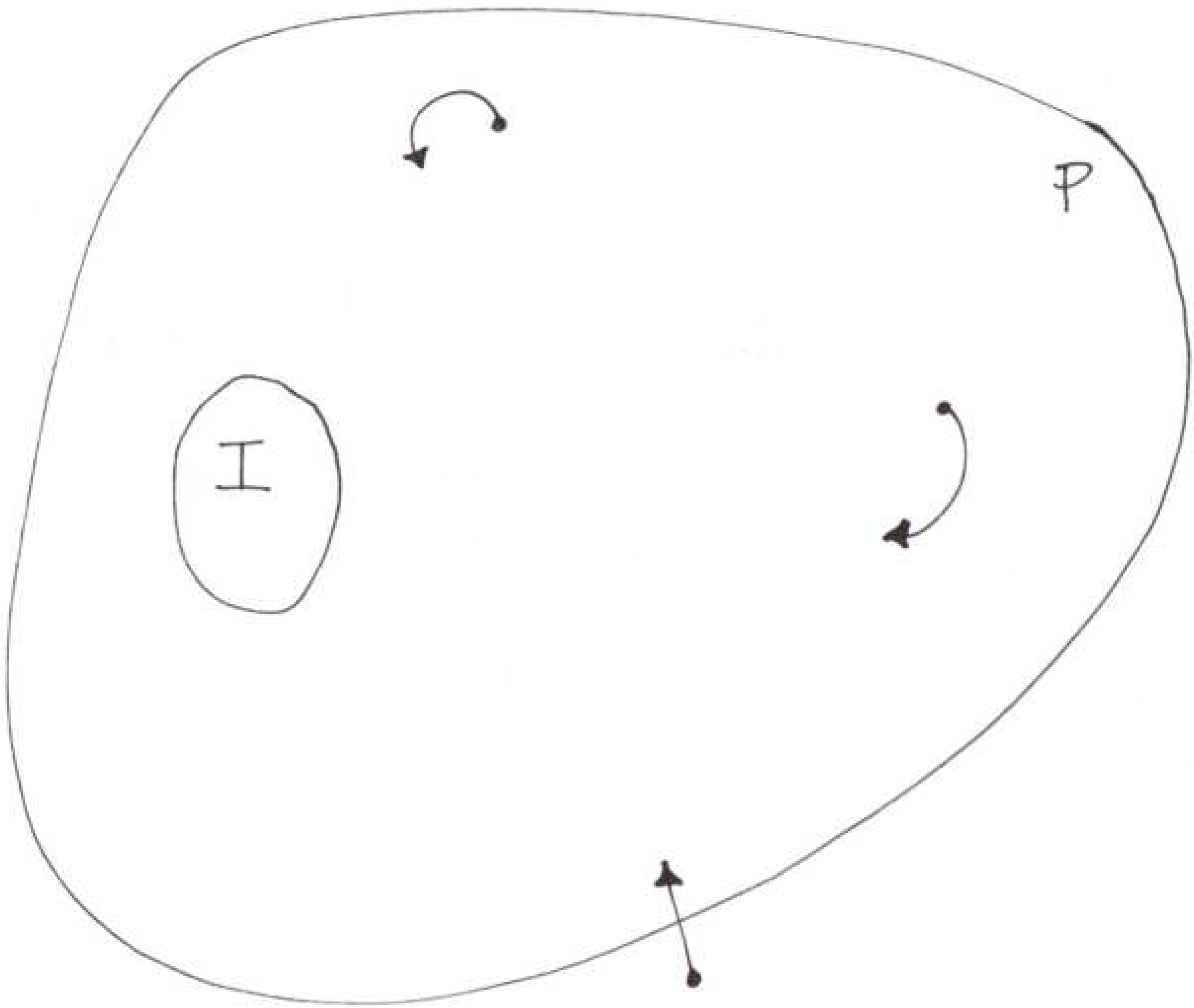
To prove that $S : (I, T)$ has safety property P, prove:

- Base case (**initiation**):

$$I \Rightarrow P$$

- Inductive case (**consecution**):

$$P \wedge T \Rightarrow P'$$

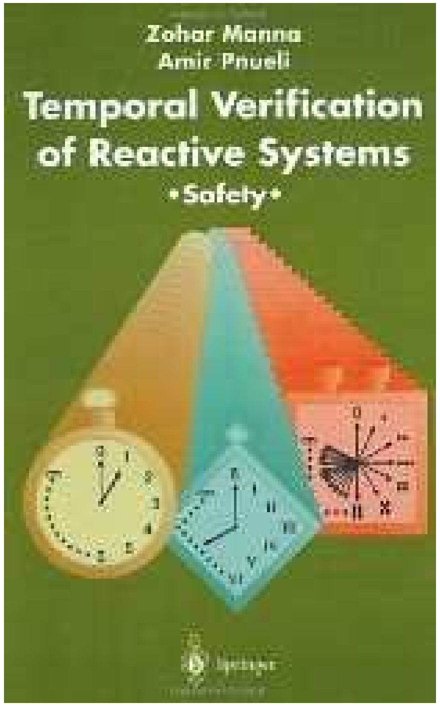


P is inductive

When Induction Fails

We present two solutions. . .

1. Use a stronger assertion, or
2. Construct an incremental proof, using previously established invariants.



– Manna and Pnueli

Temporal Verification of Reactive Systems: Safety

1995

Method 1 = “Monolithic”

Method 2 = “Incremental”

Outline

1. Illustration of the two methods
2. SAT-based model checkers
3. Understanding IC3
4. After IC3: Temporal Logics, SMT
5. Challenges for SAT/SMT

Two Transition Systems

S_1:

```
x, y := 1, 1
while *:
  x, y := x + 1, y + x
```

S_2:

```
x, y := 1, 1
while *:
  x, y := x + y, y + x
```

$$P : \; y \geq 1$$

Induction on System 1

S_1:

```
x, y := 1, 1
while *:
  x, y := x + 1, y + x
```

- Initiation:

$$\underbrace{x = 1 \land y = 1}_{\text{initial condition}} \Rightarrow \underbrace{y \geq 1}_{P}$$

- Consecution (fails):

$$\underbrace{y \geq 1}_{P} \land \underbrace{x' = x + 1 \land y' = y + x}_{\text{transition relation}} \not\Rightarrow \underbrace{y' \geq 1}_{P'}$$

Incremental Proof

S_1:

```
x, y := 1, 1
while *:
  x, y := x + 1, y + x
```

Problem: y decreases if x is negative. But...

$\varphi_1 : \; x \geq 0$

- Initiation:

$$x = 1 \wedge y = 1 \Rightarrow x \geq 0$$

- Consecution:

$$\underbrace{x \geq 0}_{\varphi_1} \wedge \underbrace{x' = x + 1 \wedge y' = y + x}_{\text{transition relation}} \Rightarrow \underbrace{x' \geq 0}_{\varphi_1'}$$

Back to P

S_1:

```
x, y := 1, 1
while *:
  x, y := x + 1, y + x
```

Consecution:

$$\underbrace{x \geq 0}_{\varphi_1} \wedge \underbrace{y \geq 1}_{P} \wedge \underbrace{x' = x + 1 \wedge y' = y + x}_{\text{transition relation}} \Rightarrow \underbrace{y' \geq 1}_{P'}$$

P is inductive **relative to** φ_1.

Induction on System 2

S_2:

```
x, y := 1, 1
while *:
  x, y := x + y, y + x
```

Induction fails for P as in System 1.
Additionally,

$$x \geq 0 \wedge x' = x + y \wedge y' = y + x \not\Rightarrow x' \geq 0$$

$x \geq 0$ is not inductive, either.

Monolithic Proof

S_2:

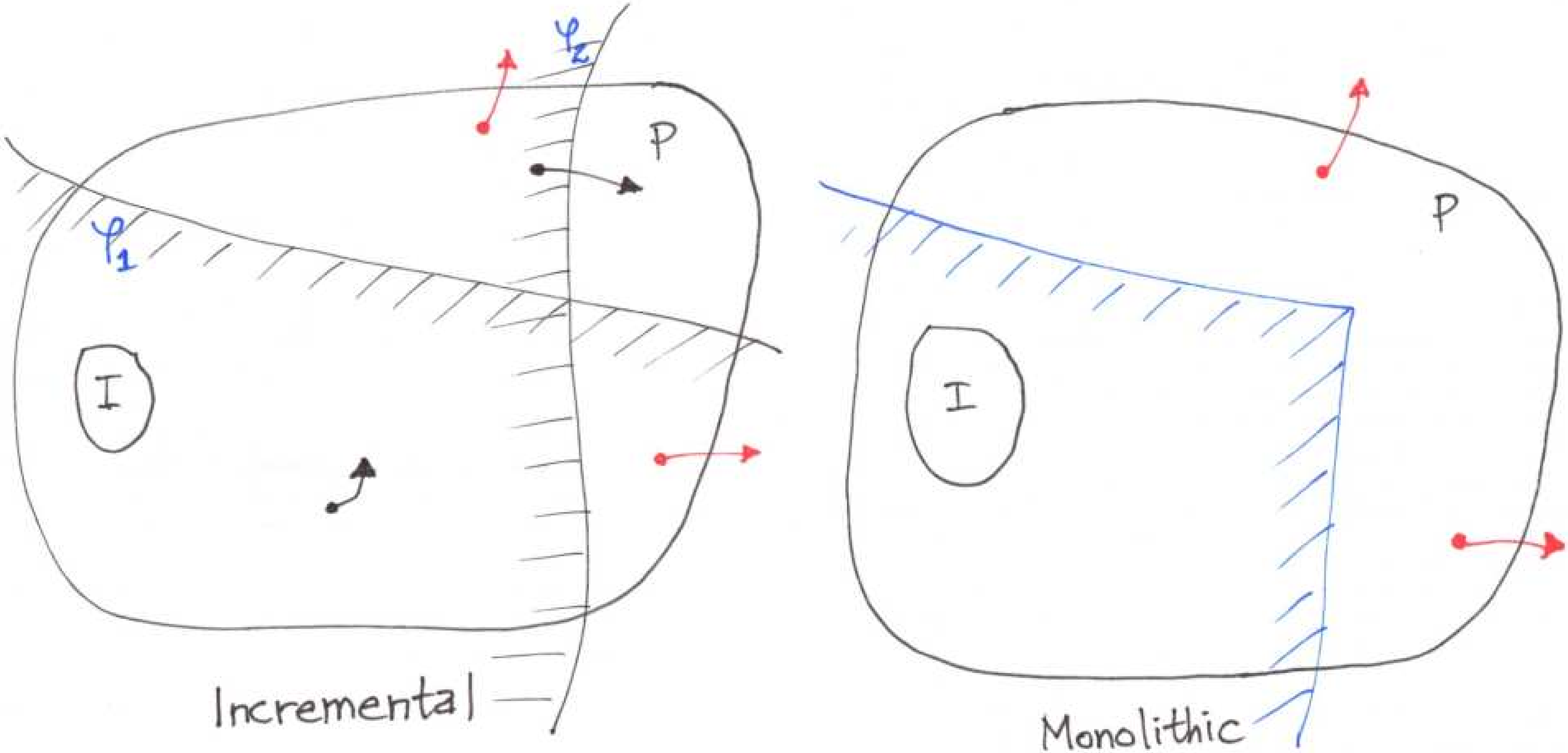
```
x, y := 1, 1
while *:
  x, y := x + y, y + x
```

Invent strengthening all at once:

$$\widehat{P}: \quad x \geq 0 \land y \geq 1$$

Consecution:

$$\underbrace{x \geq 0 \land y \geq 1}_{\widehat{P}} \land x' = x + y \land y' = y + x \Rightarrow \underbrace{x' \geq 0 \land y' \geq 1}_{\widehat{P}'}$$


φ_2
P
φ_1
I
Incremental
P
I
Monolithic

Incremental vs. Monolithic Methods

- Incremental: does not always work
- Monolithic: relatively complete
- Incremental: apply induction iteratively (“modular”)
- Monolithic: invent one strengthening formula

We strongly recommend its use whenever applicable. Its main advantage is that of **modularity**.

– Manna and Pnueli

Temporal Verification of Reactive Systems: Safety

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Finite-state System

Transition system:

$$S: \quad (\overline{i}, \overline{x}, I(\overline{x}), T(\overline{x}, \overline{i}, \overline{x}'))$$

Cube s:

- Conjunction of literals, e.g.,

$$x_1 \wedge \neg x_2 \wedge \neg x_3 \wedge x_4 \wedge \cdots$$

- Like any formula, represents set of states (that satisfy it)

Clause: $\neg s$

SAT-Based Backward Model Checking

1. Search for predecessor s to some error state:

$$P \wedge T \Rightarrow P'$$

   If none, property holds.

2. Reduce cube s to $\bar{s}$:
   - Expand to others with bad successors [McMillan 2002], [Lu et al. 2005]
   - If $P \wedge \neg s \wedge T \Rightarrow \neg s'$, reduce by implication graph [Lu et al. 2005]
   - Apply inductive generalization [Bradley 2007]

3. P := $P \wedge \neg \bar{s}$

Inductive Generalization

Given: cube s
Find: $c \subseteq \neg s$ such that

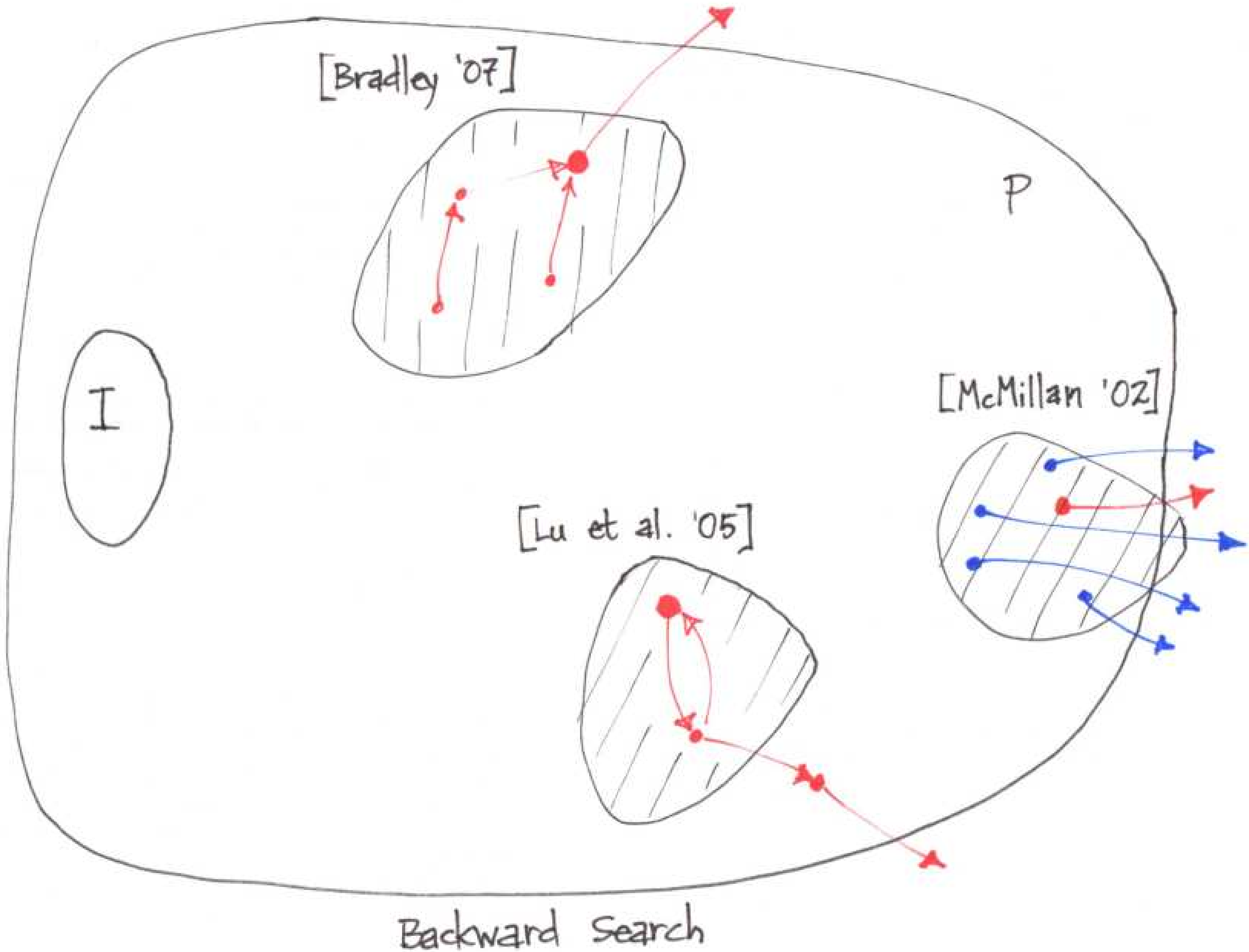
- Initiation:

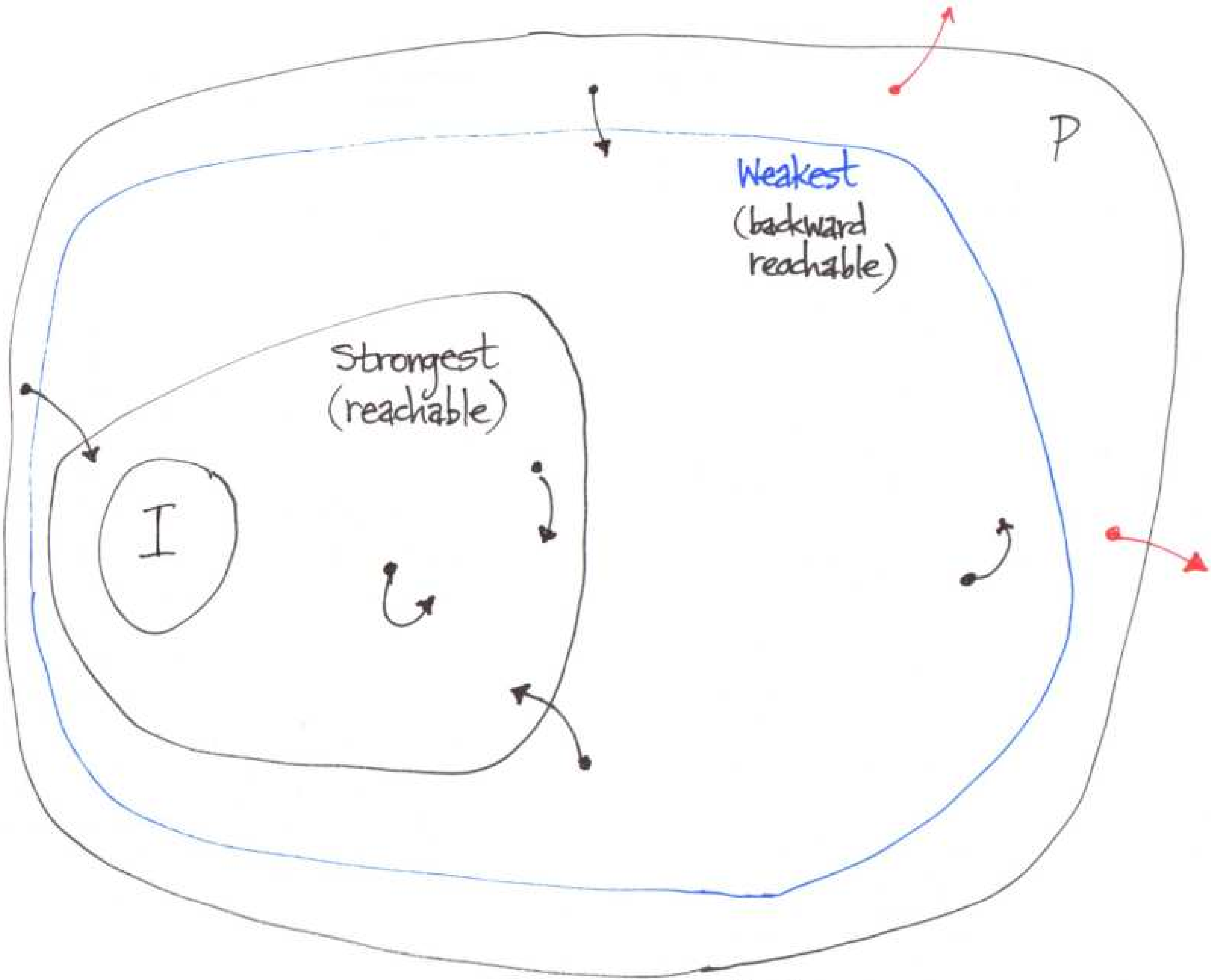
$$I \Rightarrow c$$

- Consecution (relative to information P):

$$P \wedge c \wedge T \Rightarrow c'$$

- No strict subclause of c is inductive relative to P


[Bradley '07]
P
I
[McMillan '02]
[Lu et al. '05]
Backward Search



Inductive Strengthening

Analysis of Backward Search

Strengths:

- Easy SAT queries, low memory
- Property focused
- Some are approximating, computing neither strongest nor weakest strengthening

Weaknesses:

- Essentially undirected search (bad for bug finding)
- Ignore initial states

Analysis of FSIS [Bradley 2007]

Strengths (essentially, great when it works):

- Can significantly reduce backward search
- Can find strong lemmas with induction

Weaknesses:

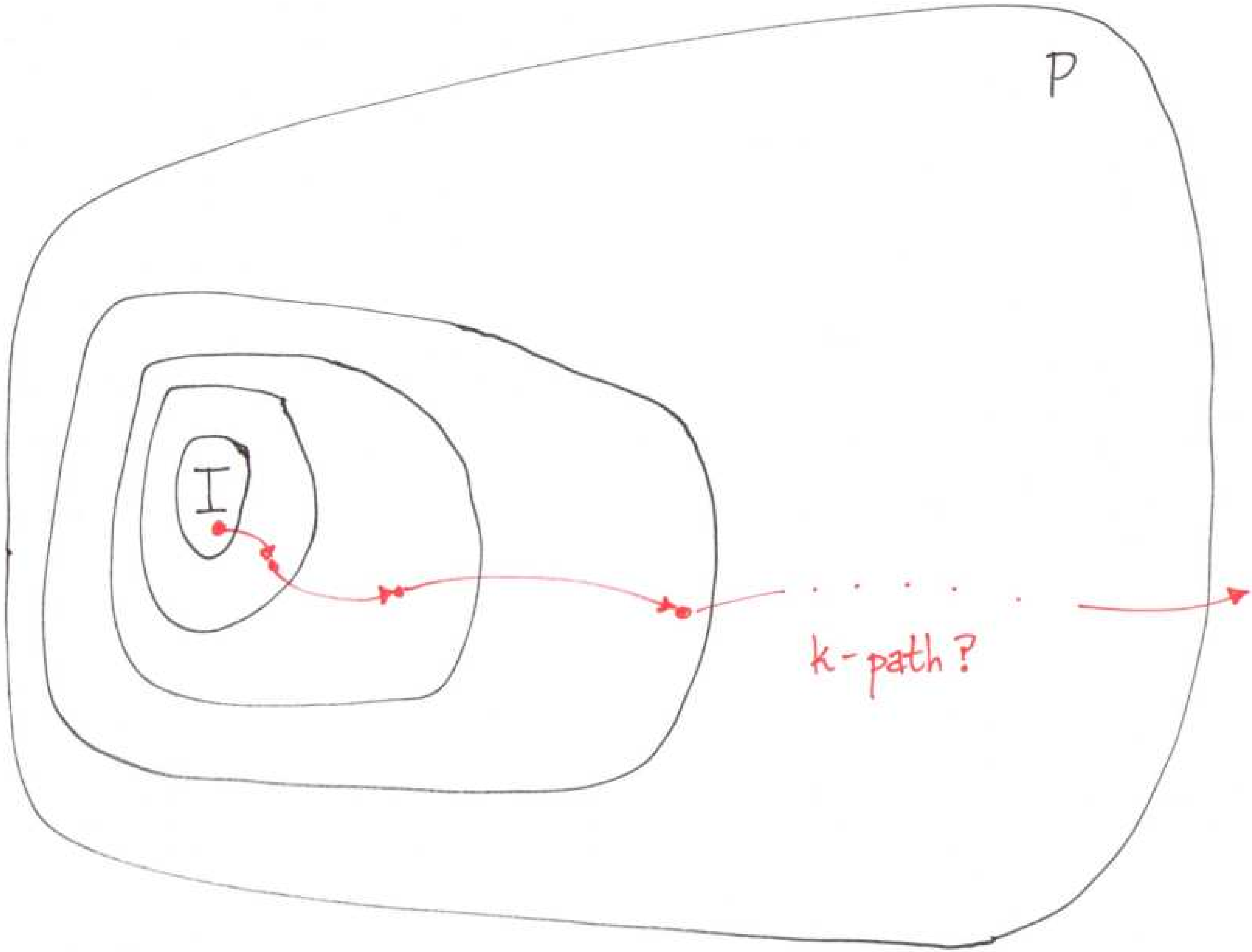
- Like others when inductive generalization fails

BMC [Biere et al. 1999]

Compared to backward search:

- Considers initial and final states
- Requires solving hard SAT queries
- Practically incomplete (UNSAT case)

$$I \wedge \bigwedge_{i=0}^{k-1} (P^{(i)} \wedge T^{(i)}) \wedge \neg P^{(k)}$$



BMC

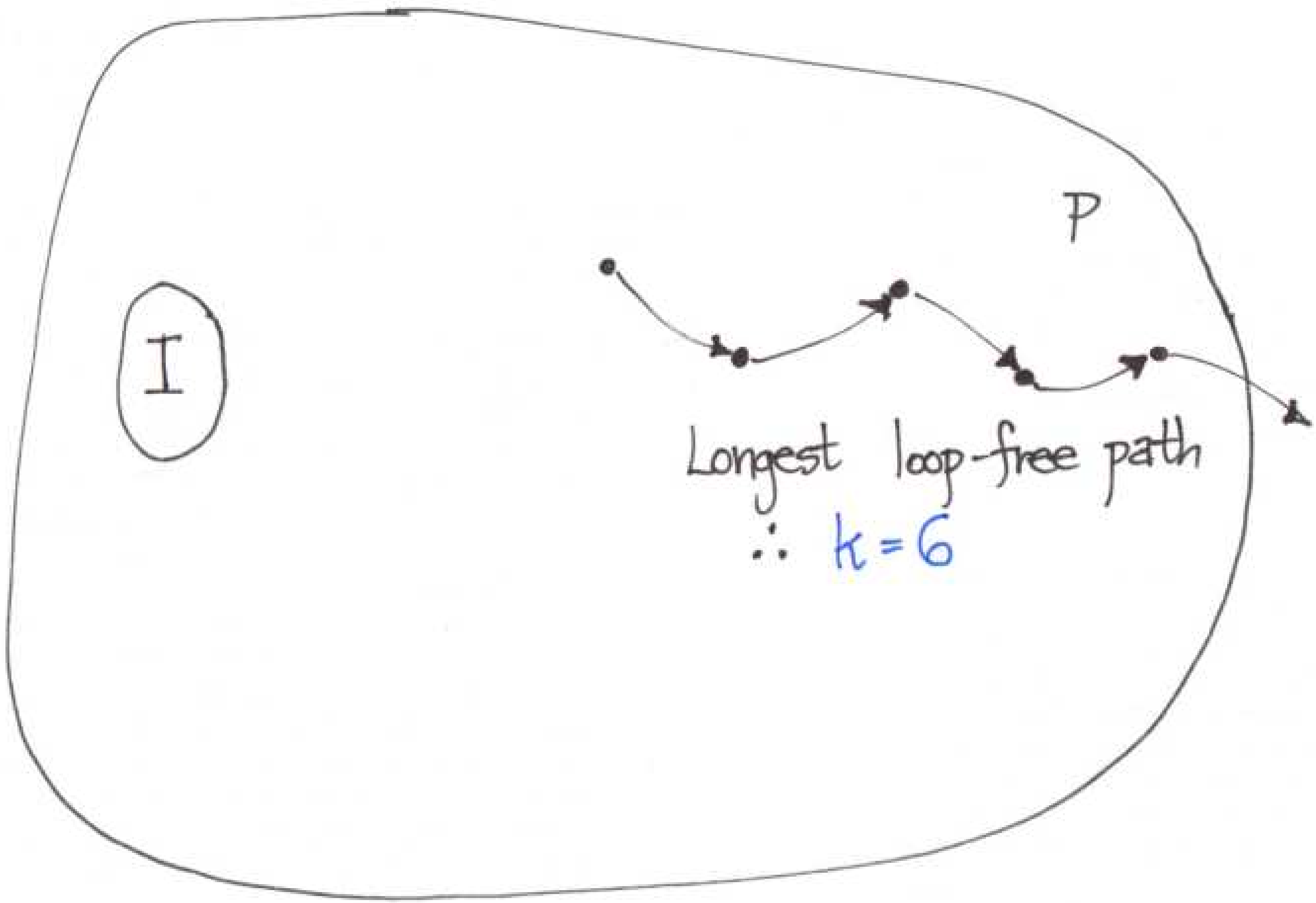
k-Induction [Sheeran et al. 2000]

Addresses practical incompleteness of BMC:

- Initiation: BMC
- Consecution:

$$\bigwedge_{i=0}^{k-1} \left(P^{(i)} \wedge T^{(i)}\right) \Rightarrow P^{(k)}$$

(plus extra constraints to consider loop-free paths)



k-Induction

ITP [McMillan 2003]

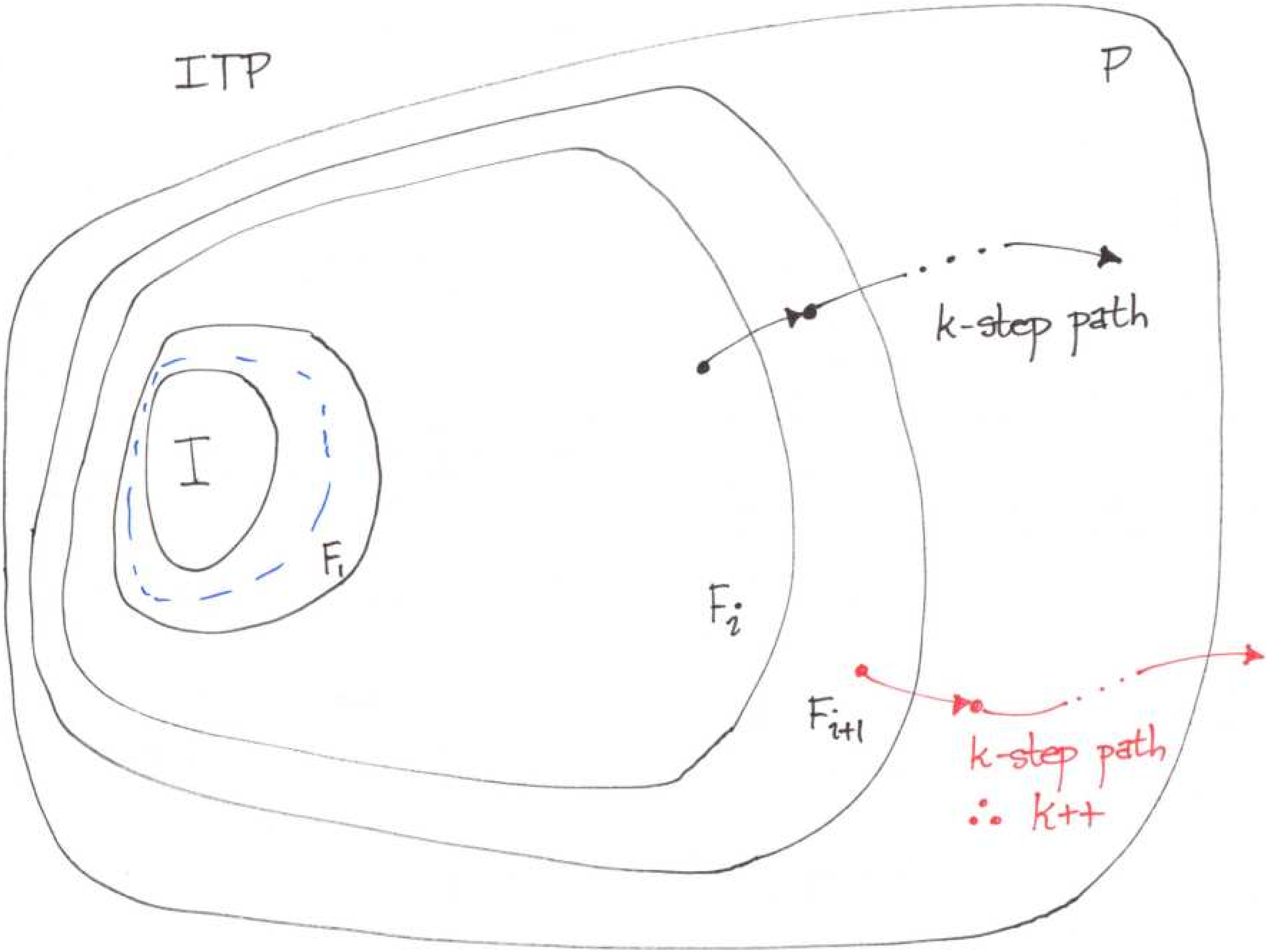
Property-focused over-approximating post-image:

$$F_i \wedge \bigwedge_{i=0}^{k-1} (P^{(i)} \wedge T^{(i)}) \Rightarrow P^{(k)}$$

- $\{$states $\leq i$ steps from initial states$\} \subseteq F_i$
- If holds, finds interpolant F_{i+1}:

$$F_i \wedge T \Rightarrow F'_{i+1} \qquad F'_{i+1} \wedge \bigwedge_{i=1}^{k-1} (P^{(i)} \wedge T^{(i)}) \Rightarrow P^{(k)}$$

- If fails, increases k


ITP
P
I
F1
Fi
Fi+1
k-step path
k-step path
∴ k++

BMC $\rightarrow$ k-Induction $\rightarrow$ ITP

- Completeness from unrolling transition relation
- Evolution: reduce max k in practice (UNSAT case)
- Monolithic:
  - hard SAT queries
  - induction at top-level only
- Consider both initial and final states

Best of Both?

Desire:

- Stable behavior (backward search)
  - Low memory, reasonable queries
  - Can just let it run
- Consideration of initial and final states (BMC)
- Modular reasoning (incremental method)

Avoid:

- Blind search (backward search)
- Queries that overwhelm the SAT solver (BMC)

IC3: A Prover

Stepwise sets $F_0, F_1, \ldots, F_k, F_{k+1}$ (CNF):

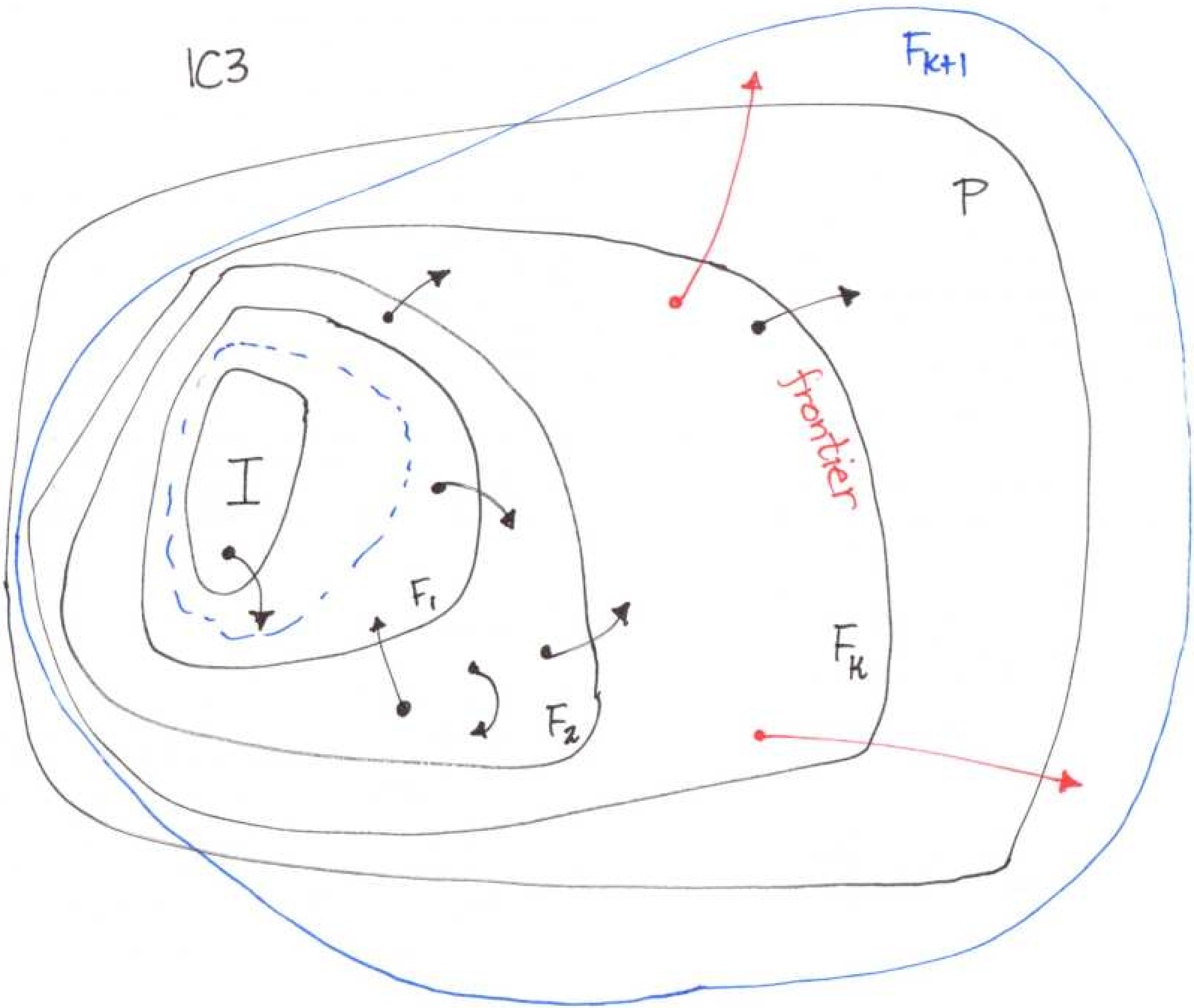
- $\{$states $\leq i$ steps from initial states$\} \subseteq F_i$
- $F_i \subseteq \{$states $\geq k - i + 1$ steps from error$\}$

Four invariants:

- $F_0 = I$
- $F_i \Rightarrow F_{i+1}$
- $F_i \wedge T \Rightarrow F'_{i+1}$
- Except F_{k+1}, $F_i \Rightarrow P$

$\therefore$ if ever $F_i = F_{i+1}$, F_i is inductive & P is invariant

IC3



Essence of IC3

- Continual refinement of over-approximating stepwise sets
  - Until one is inductive
  - Monolithic use of induction
- Generation of clauses as response to backward reachable states
  - Inductive generalization: $c \subseteq \neg s$
    (c is inductive relative to a stepwise set)
  - Incremental use of induction

Two Views of IC3

- Prover: Generates predicates from counterexamples
  - From s: state that can reach error
  - To $c \subseteq \neg s$: inductive relative to F_i
  - c proves that s is unreachable in $\leq i + 1$ steps
- Bug finder: Guided backward search
  - Stepwise sets: proximity estimate to initial state

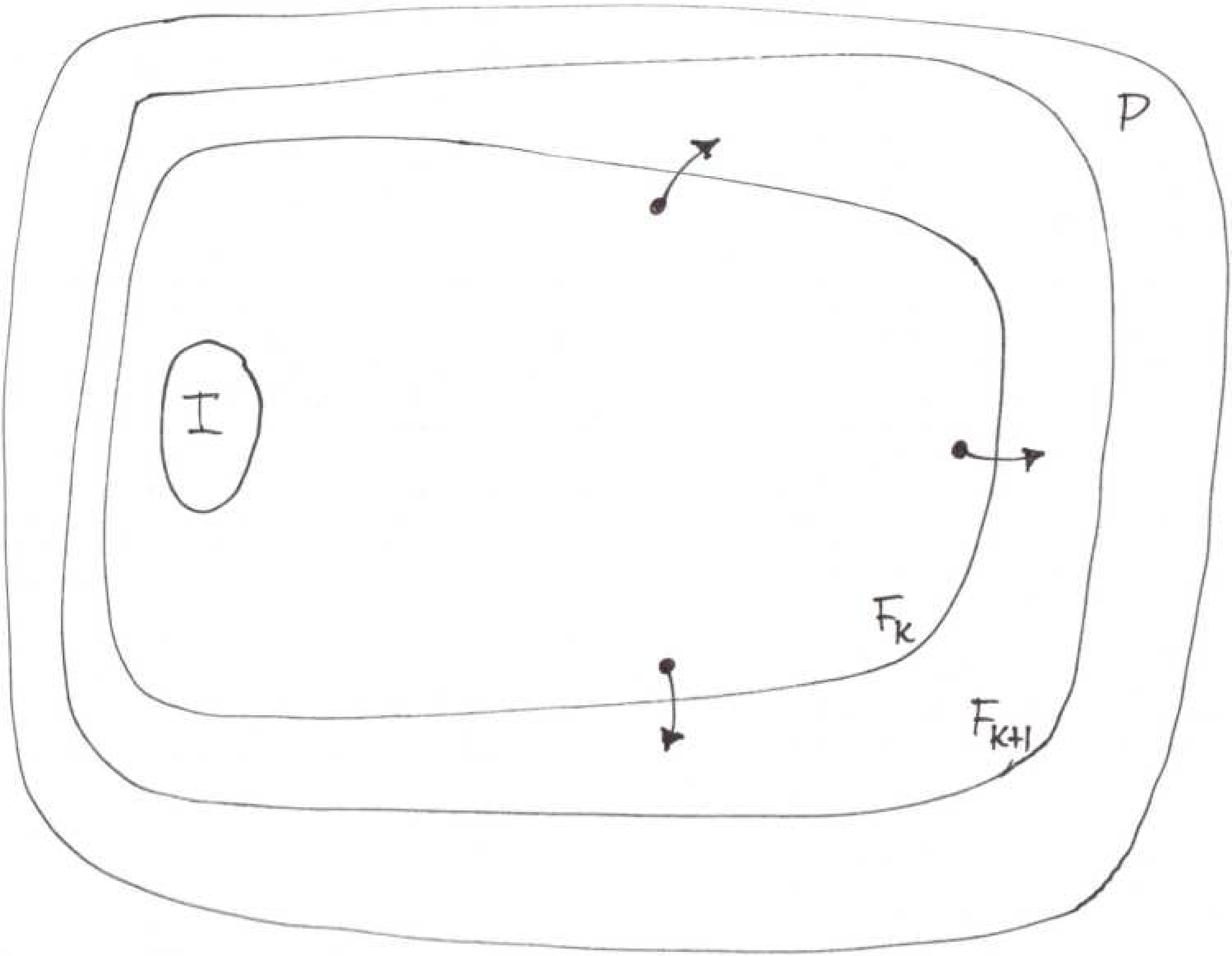
Induction at Top Level

Is P inductive relative to F_k?

$$F_k \wedge T \Rightarrow P'$$

(Recall: $F_k \Rightarrow P$)

- Possibility #1: Yes
- Conclusion: P is inductive relative to F_k



$$F_k (\land P) \land T \Rightarrow P'$$

Induction at Top Level

Monolithic behavior (predicate abstraction):

- For i from 1 to k: find largest $C \subseteq F_i$ s.t.

$$F_i \wedge T \Rightarrow C'$$

  F_{i+1} := $F_{i+1} \wedge C$

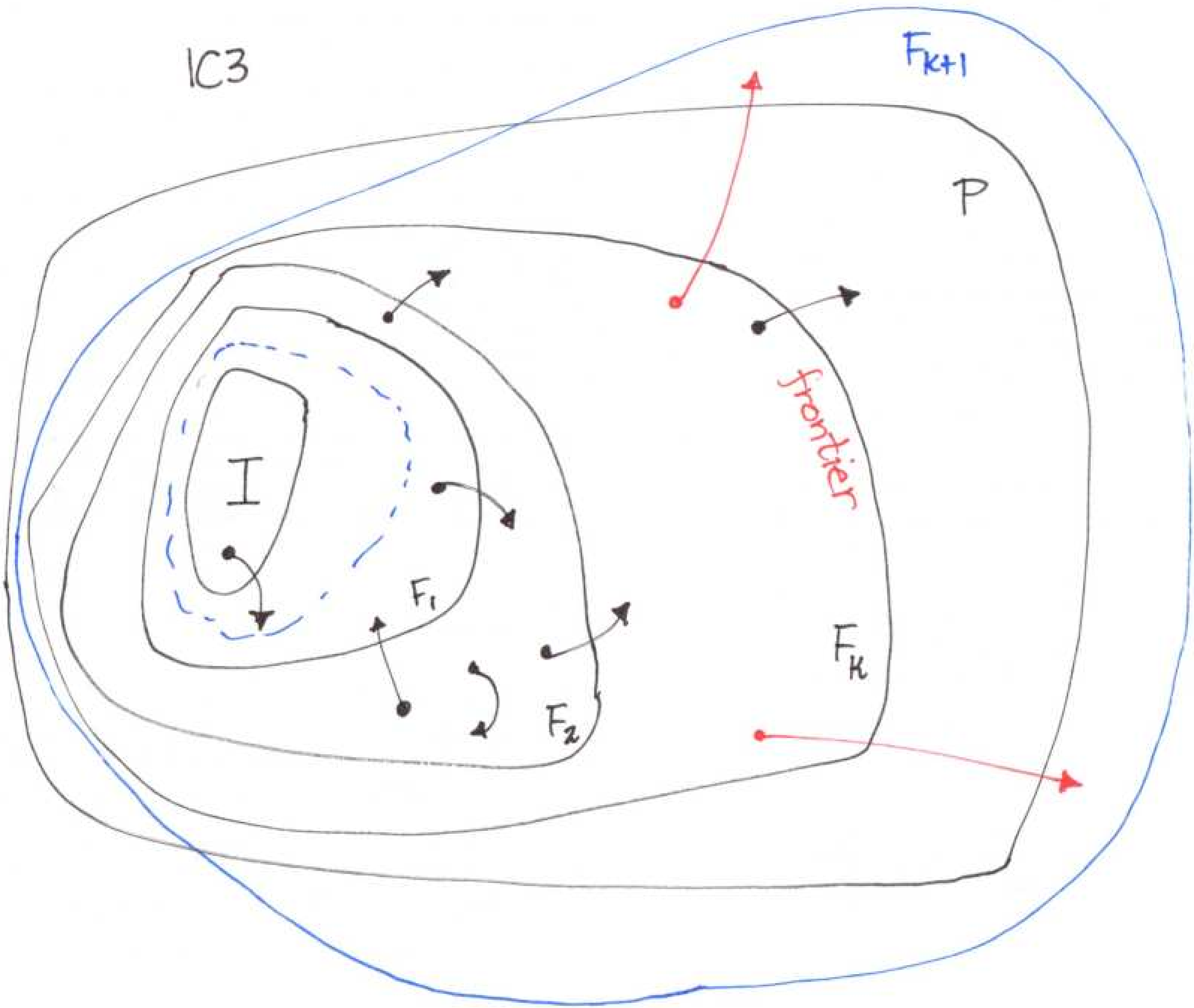
- F_{k+1} := $F_{k+1} \wedge P$
- New frontier: F_{k+1}

If ever $F_i = F_{i+1}$, done: P is invariant.

Counterexample To Induction (CTI)

$$F_k \wedge T \Rightarrow P'$$

- Possibility #2: No
- Conclusion: $\exists$ F_k-state s with error successor
- If s is an initial state, done: P is not invariant
- Otherwise...


IC3
F_{k+1}
P
frontier
I
F_1
F_2
F_k

Induction at Low Level

Inductive Generalization in IC3

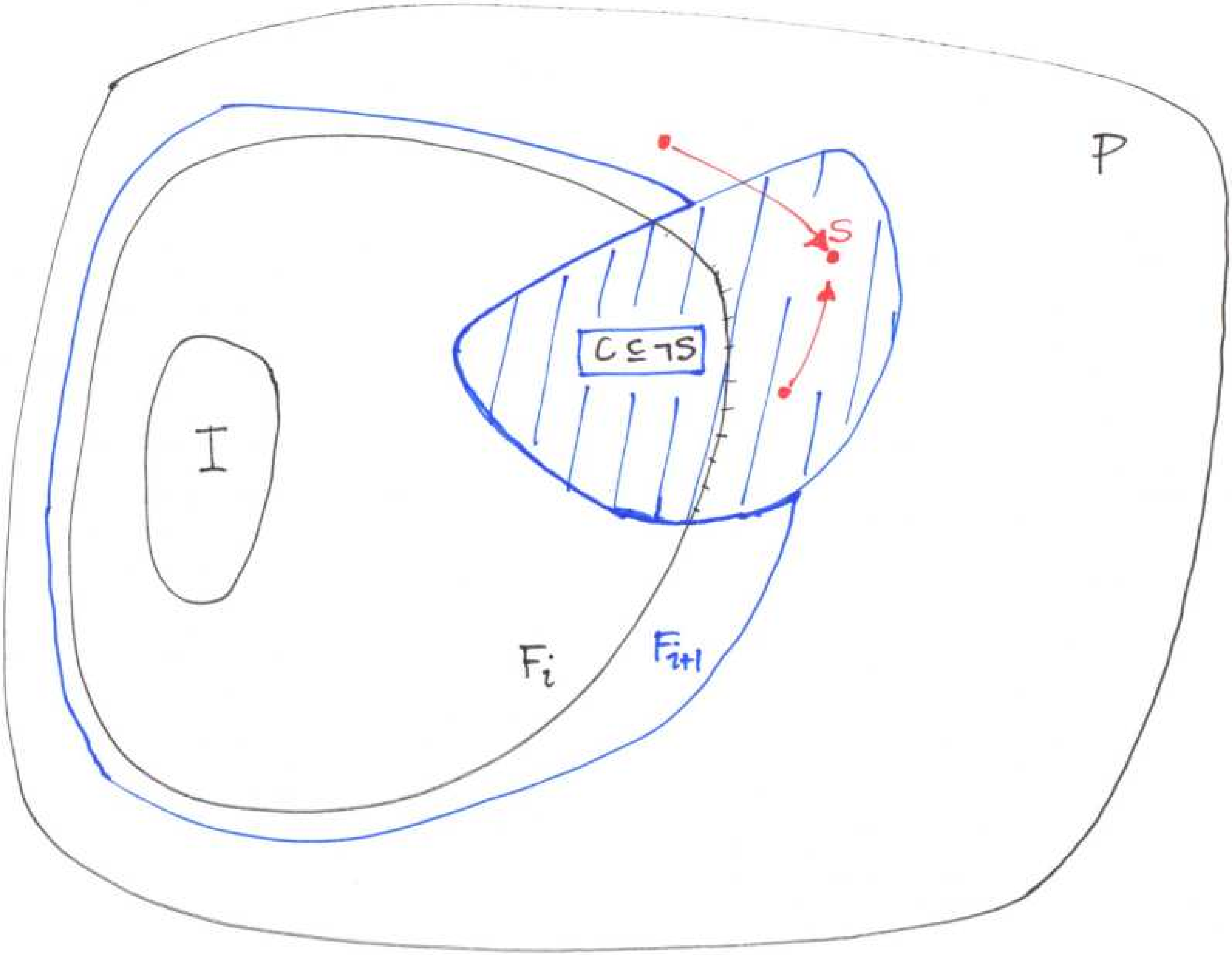
- **Given**: cube s
- **Find**: $c \subseteq \neg s$ such that
  - Initiation:

$$I \Rightarrow c$$

  - Consecution (relative to F_i):

$$F_i \wedge c \wedge T \Rightarrow c'$$

  - No strict subclause of c is inductive relative to F_i



Inductive Generalization

Addressing CTI s

- Find highest i such that

$$F_i \wedge \neg s \wedge T \Rightarrow \neg s'$$

- Apply inductive generalization:

$$c \subseteq \neg s \qquad I \Rightarrow c \qquad F_i \wedge c \wedge T \Rightarrow c'$$

- $\therefore F_{i+1}$:= $F_{i+1} \wedge c$ (also update F_j, $j \leq i$)
- If $i < k$, new **proof obligation**:

$$(s,\ i+1)$$

"Inductively generalize s relative to F_{i+1}"

Addressing Proof Obligation (t, j)

SAT query:

$$F_j \wedge \neg t \wedge T \Rightarrow \neg t'$$

If UNSAT:

- Inductive generalization must succeed:

$$c \subseteq \neg t \qquad I \Rightarrow c \qquad F_j \wedge c \wedge T \Rightarrow c'$$

- F_{j+1} := $F_{j+1} \wedge c$
- Updated proof obligation (if $j < k$): $(t,\ j+1)$

Addressing Proof Obligation (t, j)

SAT query:

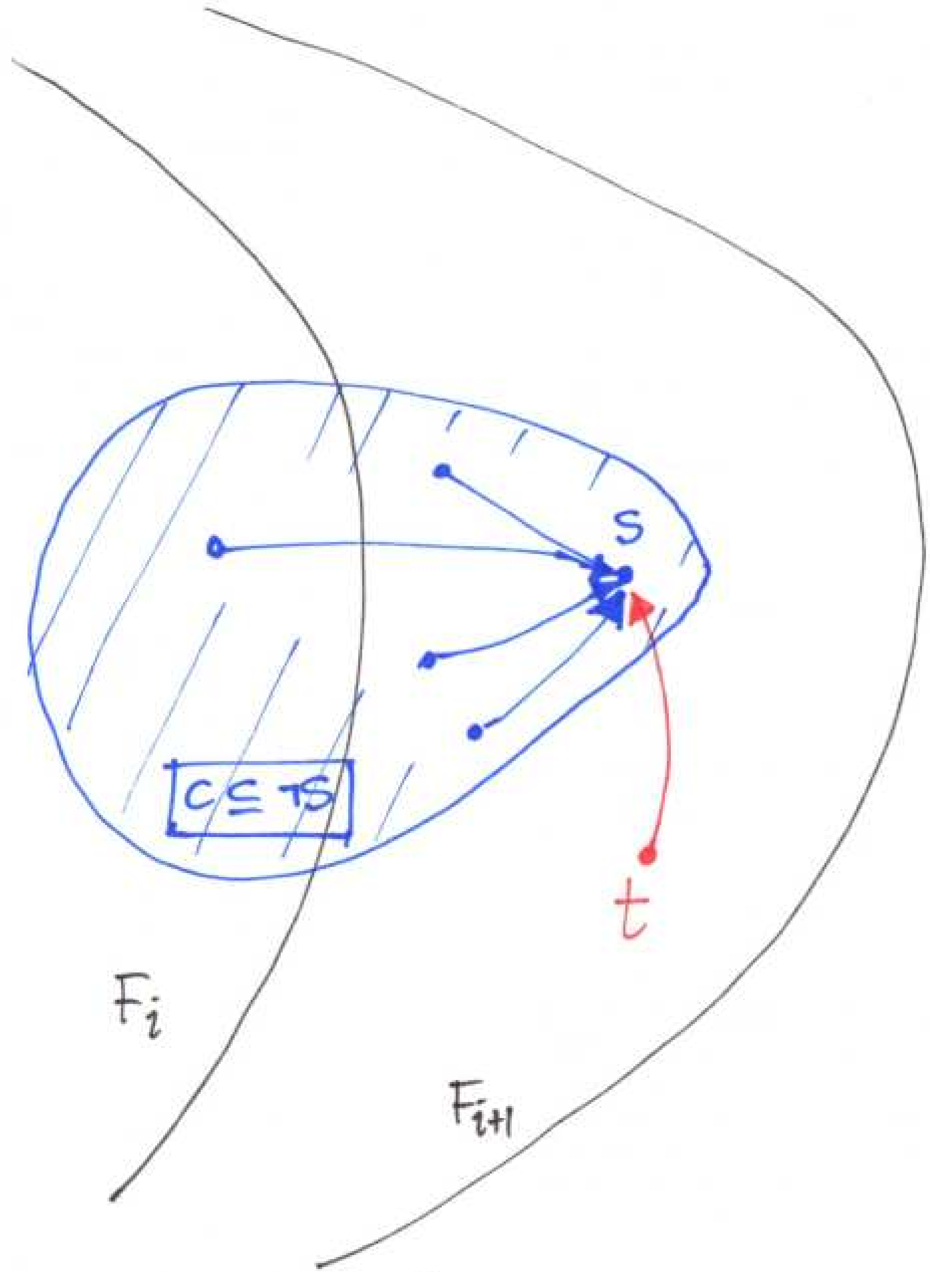
$$F_j \wedge \neg t \wedge T \Rightarrow \neg t'$$

If SAT: New CTI u, treat as before

One of IC3's Insights

Identification of **relevant predecessors**:

- Why did inductive generalization of s
  - succeed relative to F_i but
  - fail relative to F_{i+1}?
- Because of some F_{i+1}-state s-predecessor t.
- Analysis at F_i focuses IC3's choice of predecessors at F_{i+1}.



Relevant Predecessor

IC3: A Prover

- Based on CTIs (s), IC3 generates F_i-relative inductive clauses ($c \subseteq \neg s$) to refine F_i's.
- IC3 propagates clauses to prepare new frontier.
  - Some clauses may be too specific.
  - Their loss can break mutual support.
- As the frontier advances, IC3 considers ever more general situations.
- It eventually finds the real reasons (as truly inductive clauses) that P is invariant.

IC3: A Bug Finder

Suppose:

- $u \rightarrow t \rightarrow s \rightarrow$ Error
- Proof obligations:

$$\{(s,\ k-1),\ (t,\ k-2),\ (u,\ k-1)\}$$

That is,

- s must be inductively generalize relative to F_{k-1}
- t must be inductively generalize relative to F_{k-2}
- u must be inductively generalize relative to F_{k-1}

Which proof obligation should IC3 address next?

Guided Search

Two observations:

- u is the "deepest" of the states

$$u \to t \to s \to \text{Error}$$

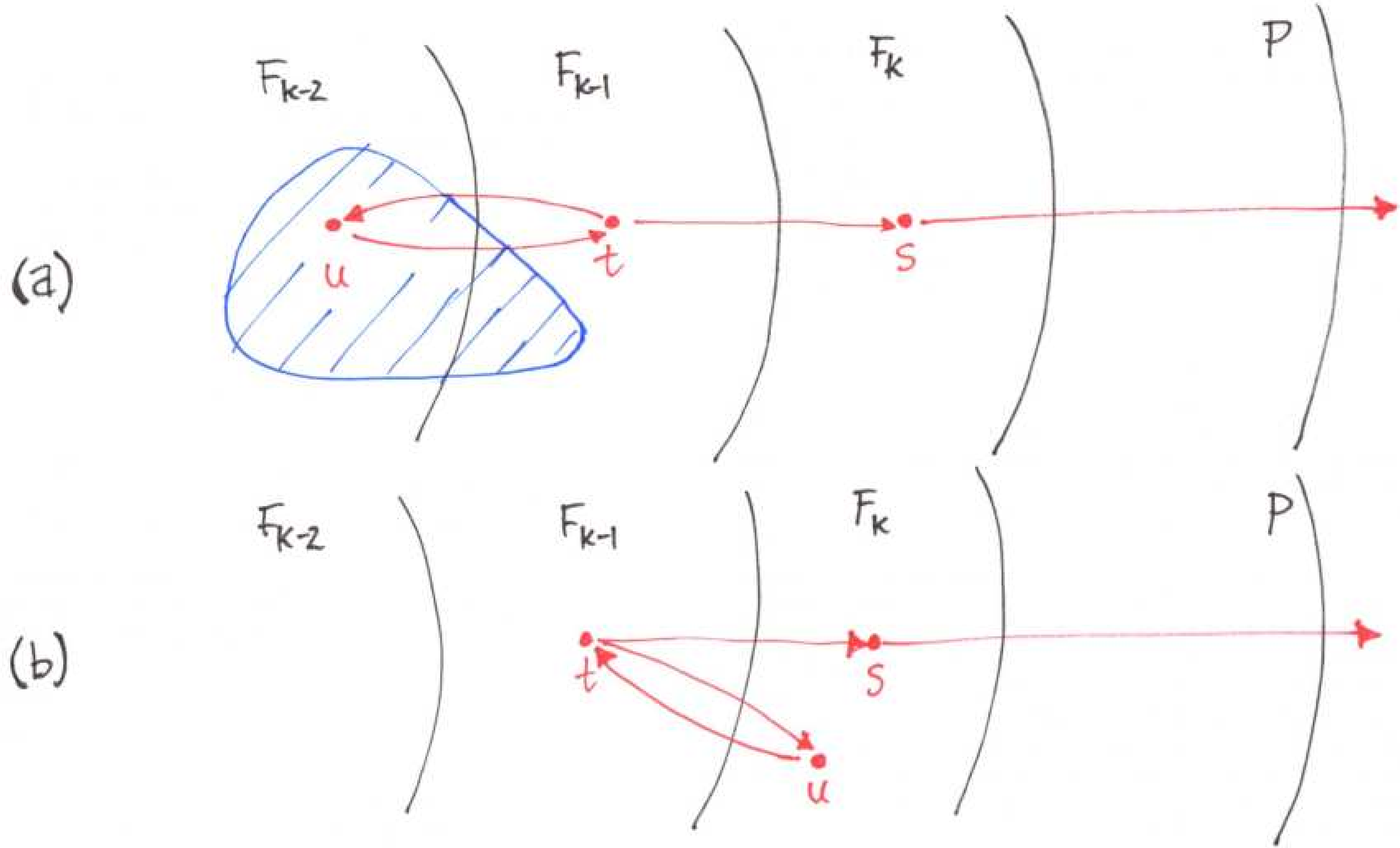
- t is the state that IC3 **considers as likeliest to be closest to an initial state**.

$$\{(s,\ k-1),\ (t,\ k-2),\ (u,\ k-1)\}$$

"Proximity metric"

Conclusion: Pursue $(t,\ k-2)$ next.

(It also happens to be the correct choice [Bradley 2011].)



$\{(s, k\text{-}1),\ (t, k\text{-}2),\ (u, k\text{-}1)\}$

Proof Obligations: Guided Search

Incremental, Inductive Verification

IIV Algorithm:

- Constructs concrete **hypotheses**
- Generates intermediate lemmas **incrementally**
- Applies **induction** many times
- **Generalizes** from hypotheses to strong lemmas

After IC3: Refinements

- New heuristic: ternary simulation cube reduction
  [Een et al., FMCAD'11]
- Industrial setting: incremental verification
  [Chockler et al., FMCAD'11]

Oh, yeah, and a name change: PDR

(Thanks, Niklas!)

PDR: Temporal Logics

- FAIR [Bradley et al., FMCAD'11]
  - For ω-regular properties, e.g., LTL
  - Insight: SCC-closed regions can be characterized inductively
- IICTL [Hassan et al., CAV'12]
  - For CTL properties
  - Insight: EX (SAT), EU (IC3), EG (FAIR)
  - Standard traversal of CTL property's parse tree
    - Over- and under-approximating sets
    - Task state-driven refinement

PDR: Infinite-state Systems

- **SMT-based Induction Methods for Timed Systems** [Kindermann et al., arXiv'12]
- **Generalized Property Directed Reachability** [Hoder et al., SAT'12]
  - Boolean push-down systems
  - Linear real arithmetic
- **Software Model Checking via IC3** [Cimatti et al., CAV'12]
  - Explicit handling of CFG
  - Applies IC3 techniques to McMillan's "Lazy Abstraction with Interpolants" [McMillan, CAV'06]

PDR: Handling Proof Obligations

Some presentations use LIFO ordering:

- Trivial correctness; easier to understand
- [Hoder et al., SAT'12], [Cimatti et al., CAV'12]
- Downside: not quite as good?
  - PSPACE-complete (finite-state), so...
  - But: fixed-length counterexamples for K
  - And: not aggressive about mutual induction

Challenges for SAT/SMT

1. Emphasizes incremental calls (100s-1000s/sec) (FAIR/IICTL: even pushes/pops sets of clauses)
2. Understand effect of solver's choices on PDR
3. Variable ordering:
   - Vital to practical performance
   - Direct core assumptions and lifting
4. IIV: proofs pushed to block (e.g., FAIR)
   - Solver should report whether proof is "useful"
5. Multi-threaded access over core constraints

Conclusions

- Attempted to explain why IC3 works:
  - As a compromise between the **incremental** and **monolithic** strategies
  - Characteristics of previous SAT-based MC
  - As a prover
  - As a bug finder
- Subsequent work: temporal logic, SMT
- Challenges for SAT/SMT