



SAT-Based Verification with IC3: Foundations and Demands

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Induction

Foundation of verification for 40+ years (Floyd, Hoare)

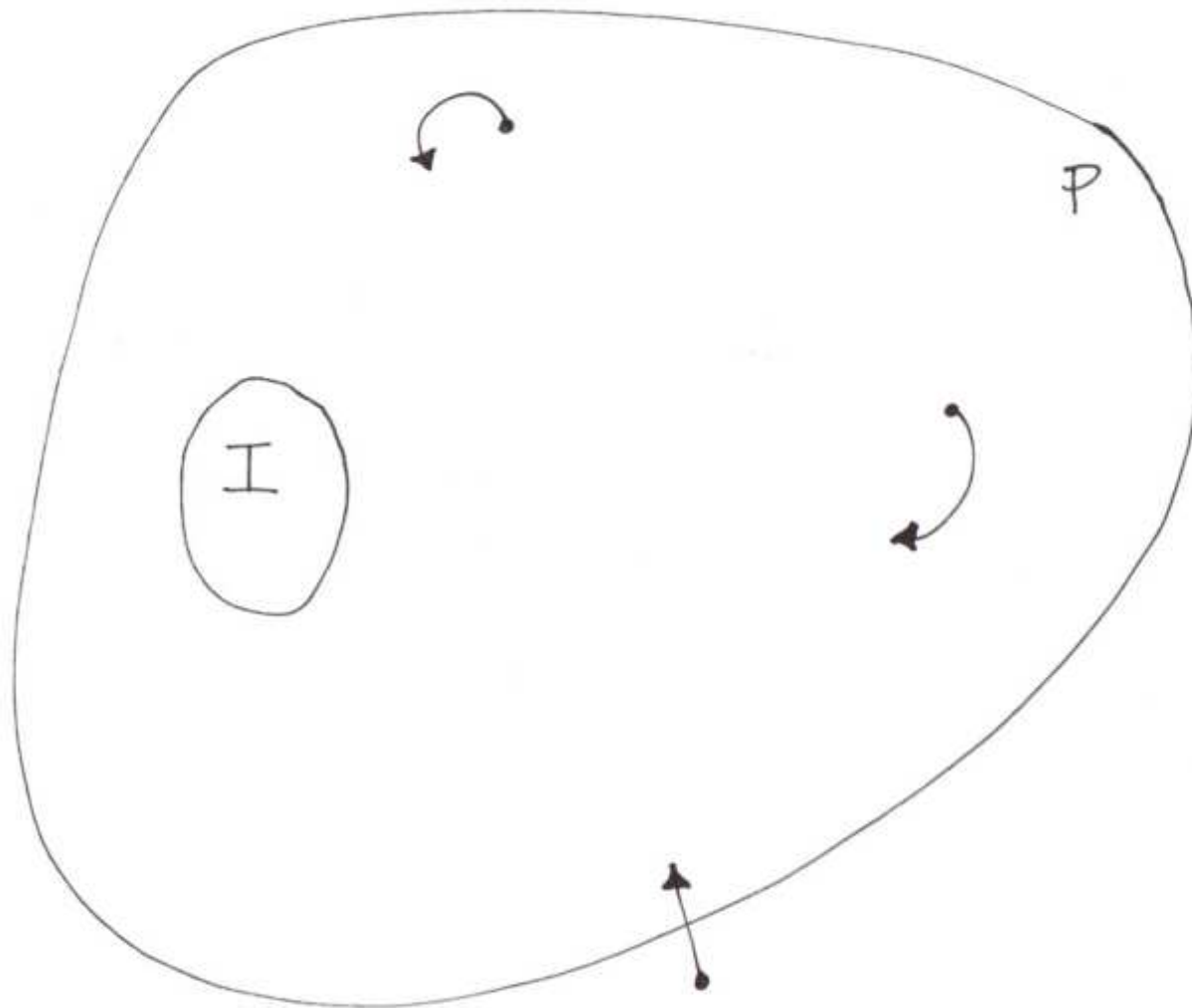
To prove that $S : (I, T)$ has safety property P , prove:

- Base case (**initiation**):

$$I \Rightarrow P$$

- Inductive case (**consecution**):

$$P \wedge T \Rightarrow P'$$

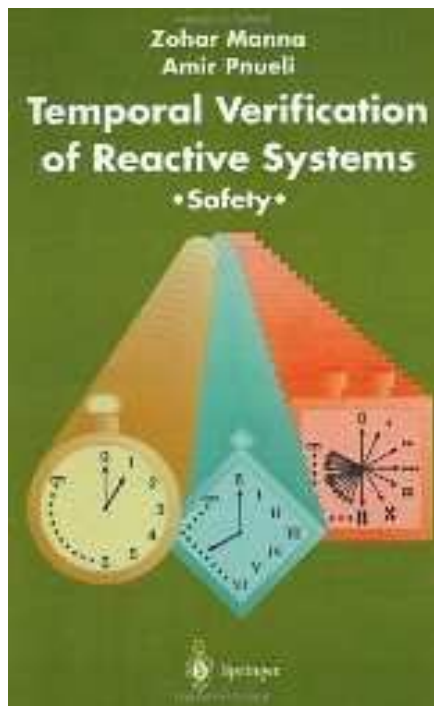


P is inductive

When Induction Fails

We present two solutions...

1. Use a stronger assertion, or
2. Construct an incremental proof, using previously established invariants.



– Manna and Pnueli
Temporal Verification of Reactive Systems: Safety
1995

Method 1 = “Monolithic”

Method 2 = “Incremental”

Outline

1. Illustration of the two methods
2. SAT-based model checkers
3. Understanding IC3
4. After IC3: Temporal Logics, SMT
5. Challenges for SAT/SMT

Two Transition Systems

S_1 :

$x, y := 1, 1$	1
while *:	2
$x, y := x + 1, y + x$	3

S_2 :

$x, y := 1, 1$	1
while *:	2
$x, y := x + y, y + x$	3

$$P : y \geq 1$$

Induction on System 1

S_1 :

$x, y := 1, 1$	1
while *:	2
$x, y := x + 1, y + x$	3

- Initiation:

$$\underbrace{x = 1 \wedge y = 1}_{\text{initial condition}} \Rightarrow \underbrace{y \geq 1}_P$$

- Consecution (fails):

$$\underbrace{y \geq 1}_P \wedge \underbrace{x' = x + 1 \wedge y' = y + x}_{\text{transition relation}} \not\Rightarrow \underbrace{y' \geq 1}_{P'}$$

Incremental Proof

S_1 :

$x, y := 1, 1$	1
while $*$:	2
$x, y := x + 1, y + x$	3

Problem: y decreases if x is negative. But...

$$\varphi_1 : x \geq 0$$

- Initiation:

$$x = 1 \wedge y = 1 \Rightarrow x \geq 0$$

- Consecution:

$$\underbrace{x \geq 0}_{\varphi_1} \wedge \underbrace{x' = x + 1 \wedge y' = y + x}_{\text{transition relation}} \Rightarrow \underbrace{x' \geq 0}_{\varphi_1}$$

Back to P

S_1 :

$x, y := 1, 1$	1
while $*$:	2
$x, y := x + 1, y + x$	3

Consecution:

$$\underbrace{x \geq 0}_{\varphi_1} \wedge \underbrace{y \geq 1}_P \wedge \underbrace{x' = x + 1 \wedge y' = y + x}_{\text{transition relation}} \Rightarrow \underbrace{y' \geq 1}_{P'}$$

P is inductive **relative to** φ_1 .

Induction on System 2

S_2 :

$x, y := 1, 1$	1
while *:	2
$x, y := x + y, y + x$	3

Induction fails for P as in System 1.
Additionally,

$$x \geq 0 \wedge x' = x + y \wedge y' = y + x \not\Rightarrow x' \geq 0$$

$x \geq 0$ is not inductive, either.

Monolithic Proof

S_2 :

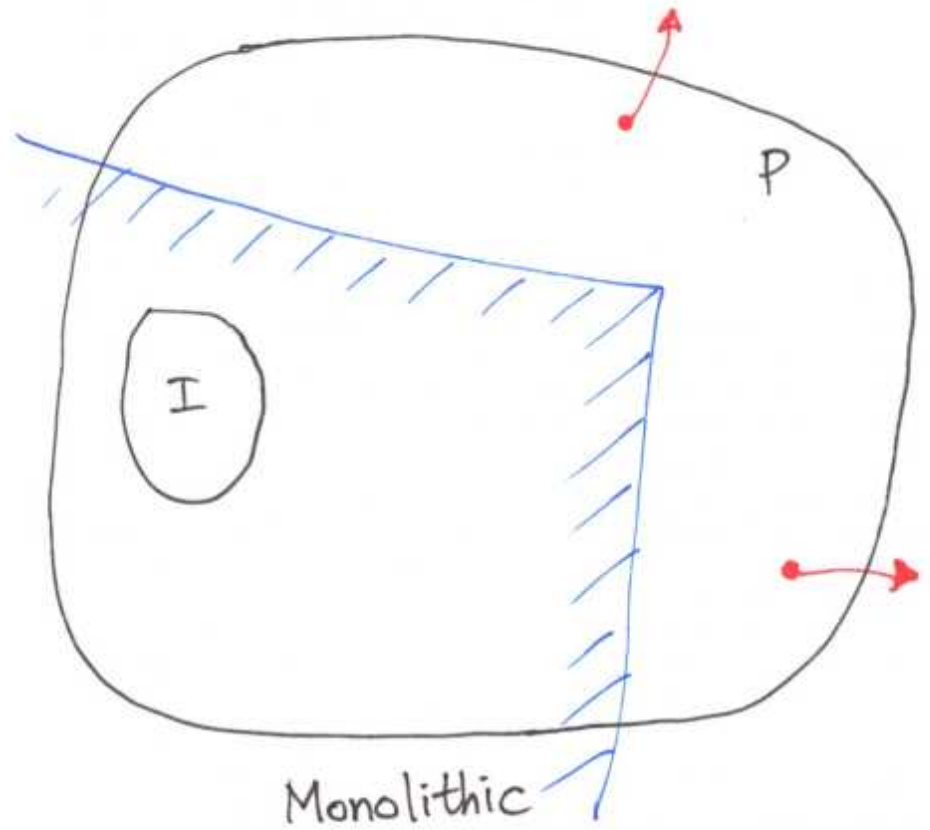
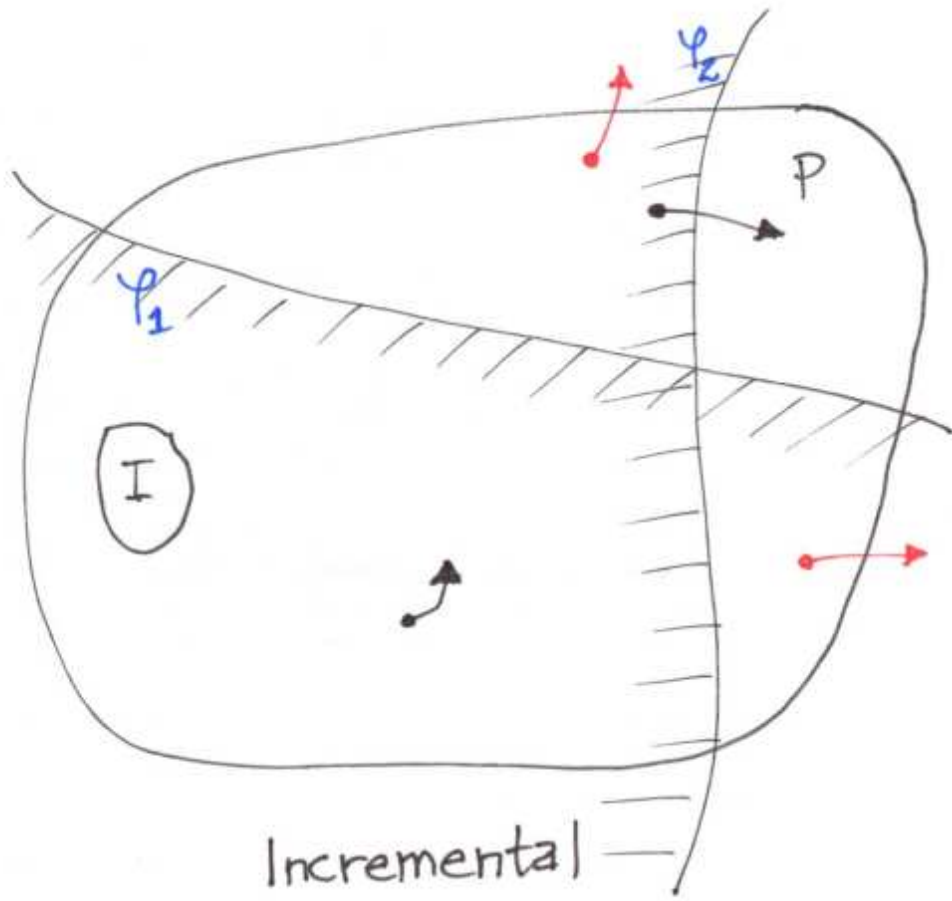
$x, y := 1, 1$	1
while *:	2
$x, y := x + y, y + x$	3

Invent strengthening all at once:

$$\hat{P} : x \geq 0 \wedge y \geq 1$$

Consecution:

$$\underbrace{x \geq 0 \wedge y \geq 1}_{\hat{P}} \wedge x' = x + y \wedge y' = y + x \Rightarrow \underbrace{x' \geq 0 \wedge y' \geq 1}_{\hat{P}'}$$



Incremental vs. Monolithic Methods

- Incremental: does not always work
- Monolithic: relatively complete
- Incremental: apply induction iteratively (“modular”)
- Monolithic: invent one strengthening formula

We strongly recommend its use whenever applicable. Its main advantage is that of **modularity**.

– Manna and Pnueli

Temporal Verification of Reactive Systems: Safety
1995

Finite-state System

Transition system:

$$S : (\bar{i}, \bar{x}, I(\bar{x}), T(\bar{x}, \bar{i}, \bar{x}'))$$

Cube s :

- Conjunction of literals, e.g.,

$$x_1 \wedge \neg x_2 \wedge \neg x_3 \wedge x_4 \wedge \dots$$

- Like any formula, represents set of states (that satisfy it)

Clause: $\neg s$

SAT-Based Backward Model Checking:

1. Search for predecessor s to some error state:

$$P \wedge T \Rightarrow P'$$

If none, property holds.

2. Reduce cube s to \bar{s} :

- Expand to others with bad successors
[McMillan 2002], [Lu et al. 2005]
- If $P \wedge \neg s \wedge T \Rightarrow \neg s'$, reduce by implication graph [Lu et al. 2005]
- Apply inductive generalization [Bradley 2007]

3. $P := P \wedge \neg \bar{s}$

Inductive Generalization

Given: cube s

Find: $c \subseteq \neg s$ such that

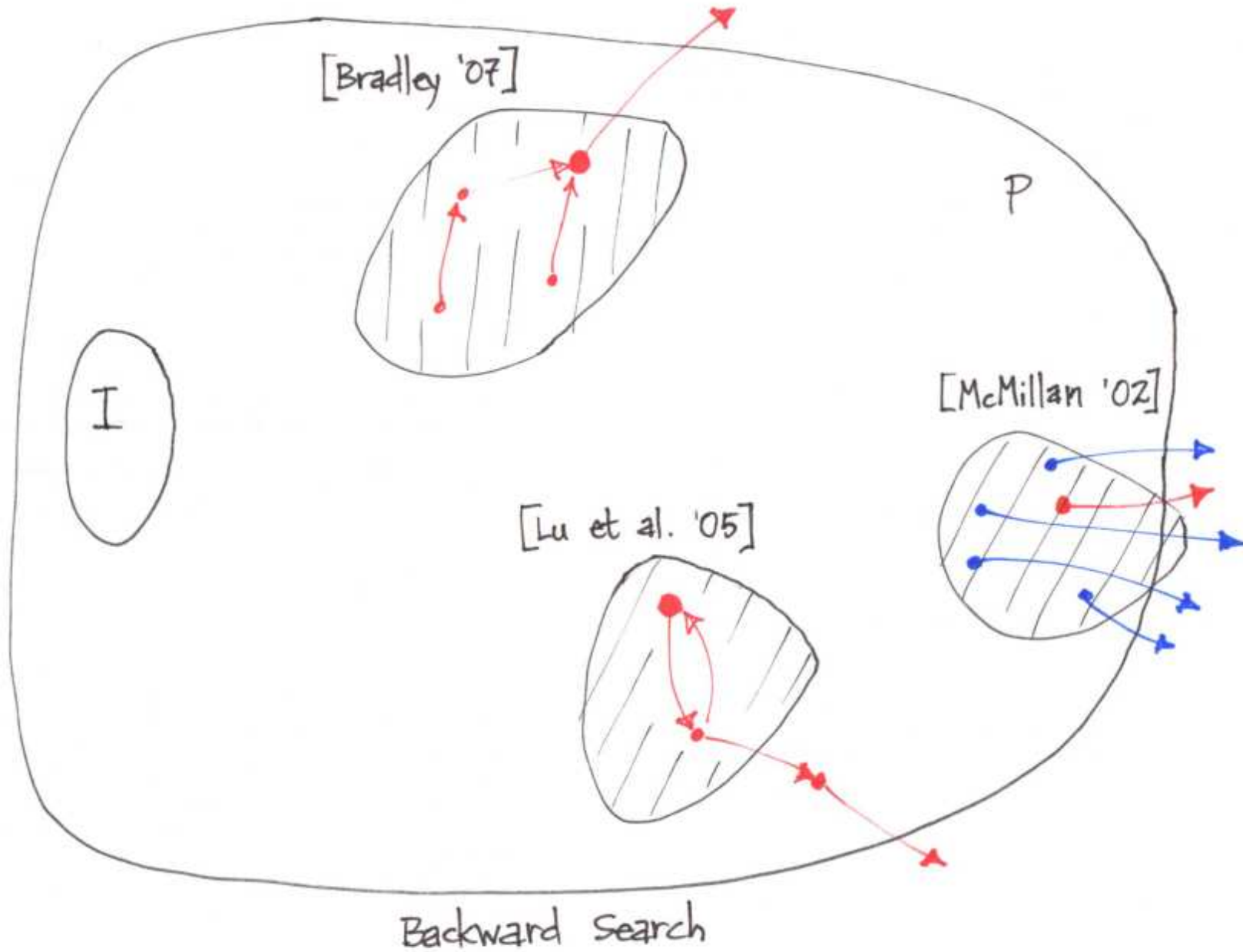
- Initiation:

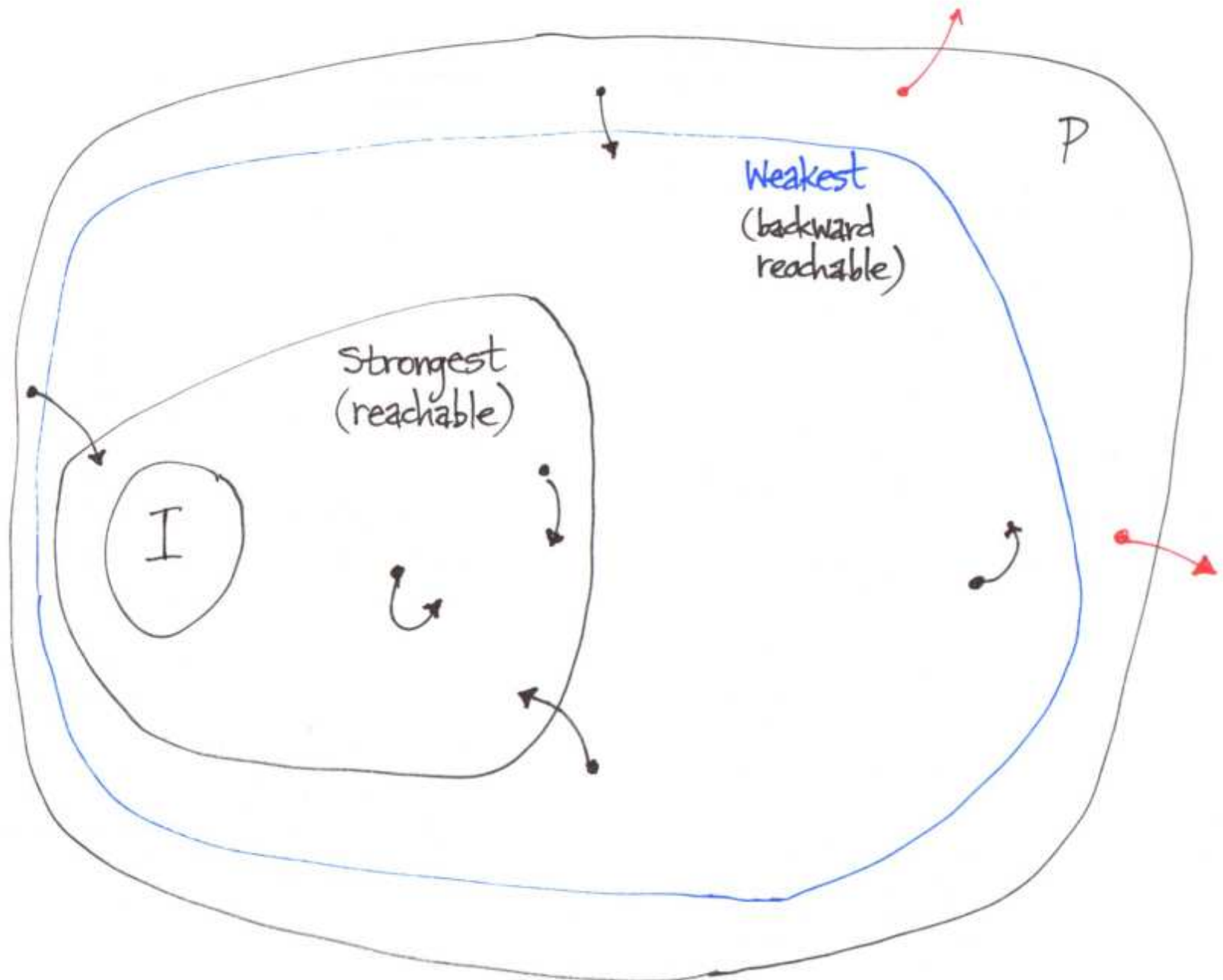
$$I \Rightarrow c$$

- Consecution (relative to information P):

$$P \wedge c \wedge T \Rightarrow c'$$

- No strict subclause of c is inductive relative to P





Inductive Strengthening

Analysis of Backward Search

Strengths:

- Easy SAT queries, low memory
- Property focused
- Some are approximating, computing neither strongest nor weakest strengthening

Weaknesses:

- Essentially undirected search (bad for bug finding)
- Ignore initial states

Analysis of FSIS [Bradley 2007]

Strengths (essentially, great when it works):

- Can significantly reduce backward search
- Can find strong lemmas with induction

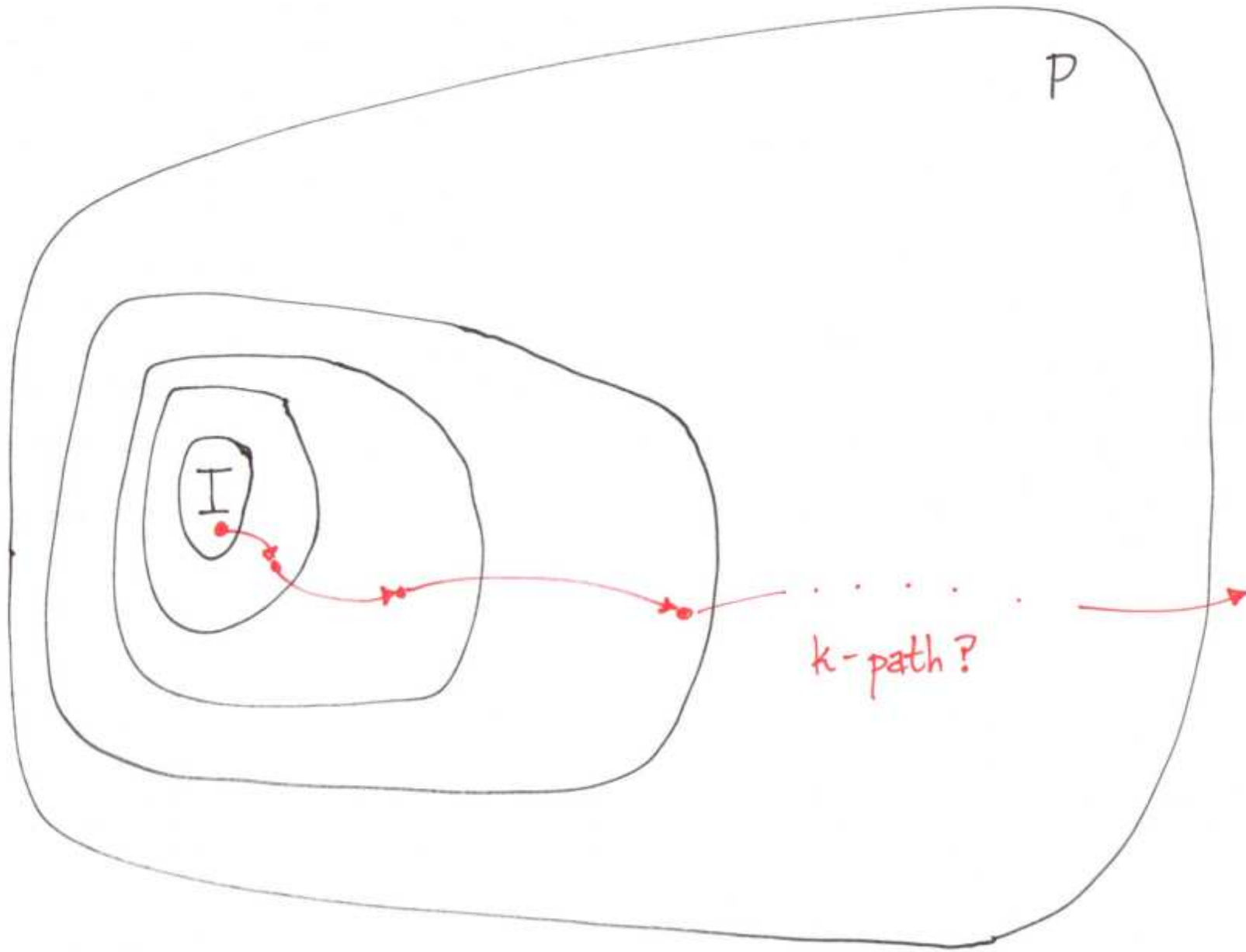
Weaknesses:

- Like others when inductive generalization fails

Compared to backward search:

- Considers initial and final states
- Requires solving hard SAT queries
- Practically incomplete (UNSAT case)

$$I \wedge \bigwedge_{i=0}^{k-1} (P^{(i)} \wedge T^{(i)}) \wedge \neg P^{(k)}$$



BMC

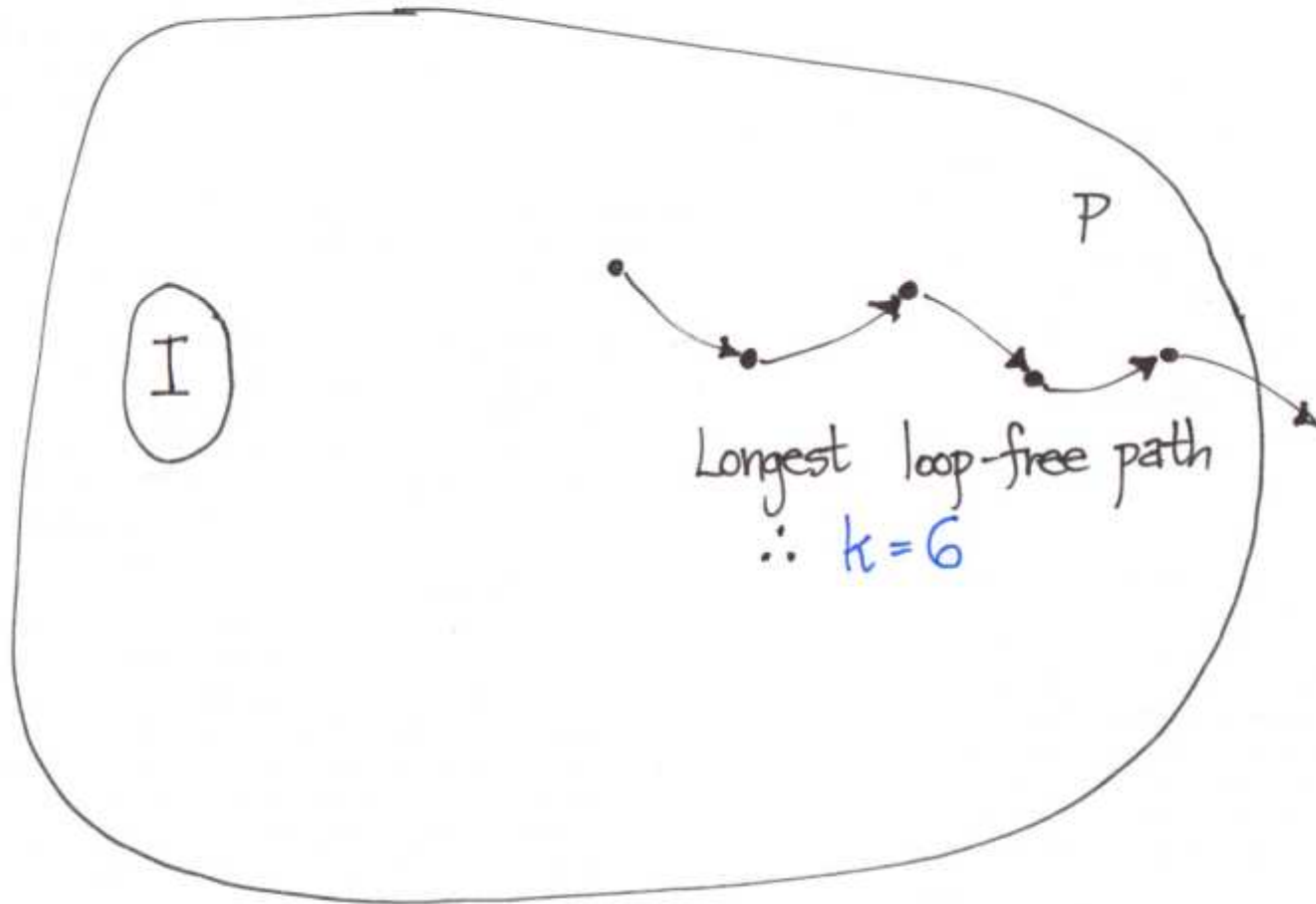
k -Induction [Sheeran et al. 2000]

Addresses practical incompleteness of BMC:

- Initiation: BMC
- Consecution:

$$\bigwedge_{i=0}^{k-1} (P^{(i)} \wedge T^{(i)}) \Rightarrow P^{(k)}$$

(plus extra constraints to consider loop-free paths)



k-Induction

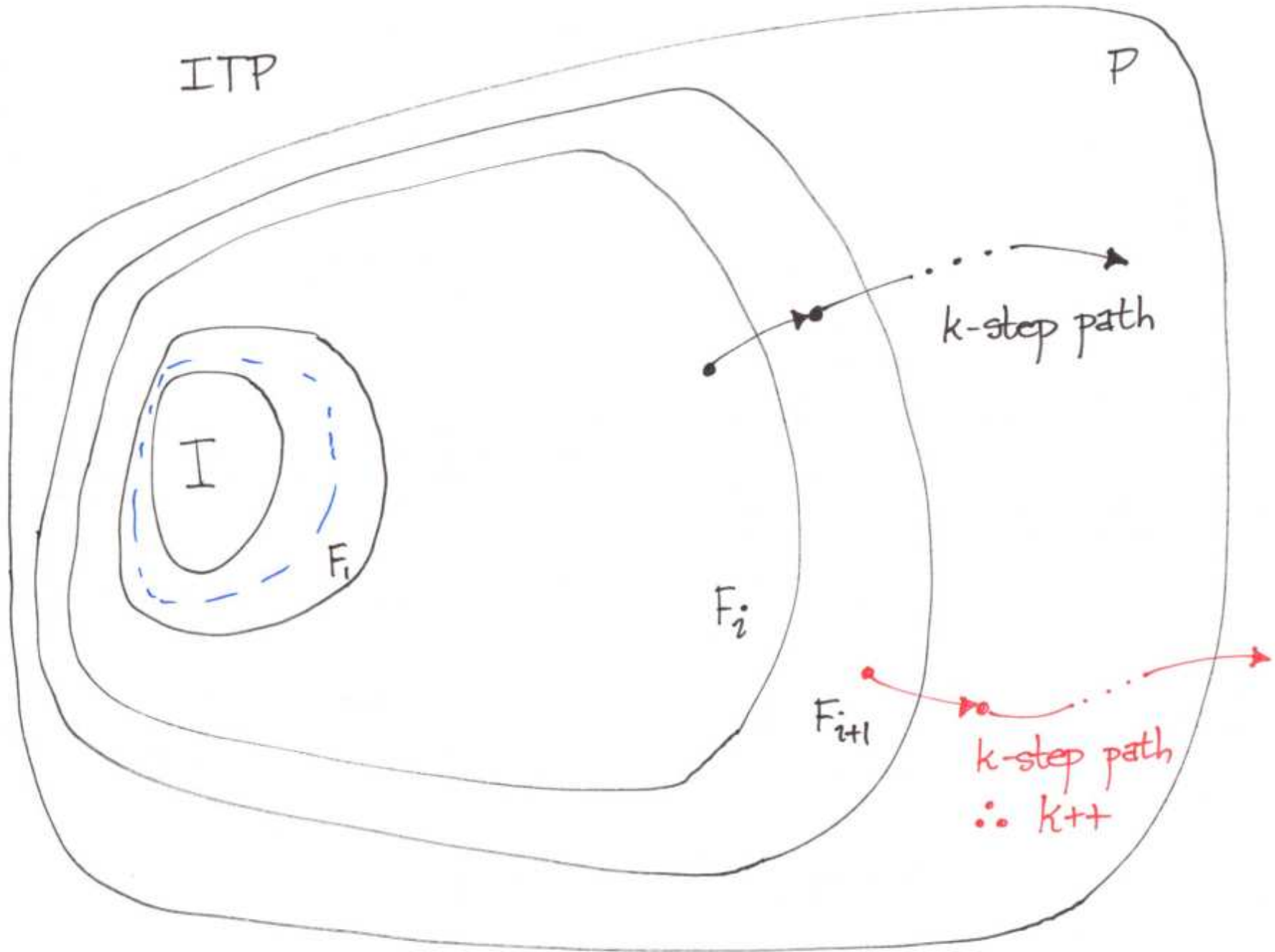
Property-focused over-approximating post-image:

$$F_i \wedge \bigwedge_{i=0}^{k-1} (P^{(i)} \wedge T^{(i)}) \Rightarrow P^{(k)}$$

- {states $\leq i$ steps from initial states} $\subseteq F_i$
- If holds, finds interpolant F_{i+1} :

$$F_i \wedge T \Rightarrow F'_{i+1} \quad F'_{i+1} \wedge \bigwedge_{i=1}^{k-1} (P^{(i)} \wedge T^{(i)}) \Rightarrow P^{(k)}$$

- If fails, increases k



BMC \rightarrow k -Induction \rightarrow ITP

- Completeness from unrolling transition relation
- Evolution: reduce max k in practice (UNSAT case)
- Monolithic:
 - hard SAT queries
 - induction at top-level only
- Consider both initial and final states

Best of Both?

Desire:

- Stable behavior (backward search)
 - Low memory, reasonable queries
 - Can just let it run
- Consideration of initial and final states (BMC)
- Modular reasoning (incremental method)

Avoid:

- Blind search (backward search)
- Queries that overwhelm the SAT solver (BMC)

IC3: A Prover

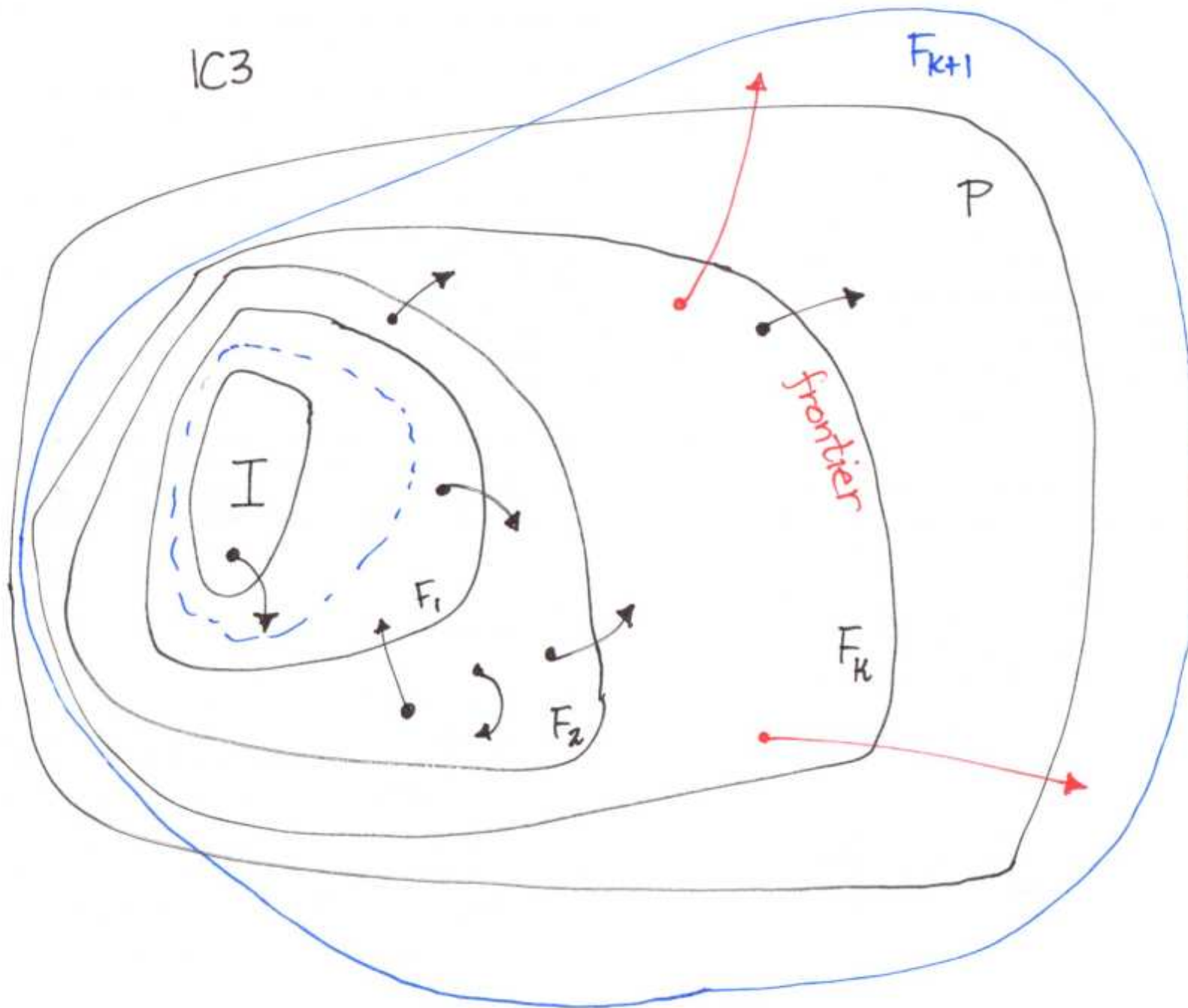
Stepwise sets $F_0, F_1, \dots, F_k, F_{k+1}$ (CNF):

- $\{\text{states} \leq i \text{ steps from initial states}\} \subseteq F_i$
- $F_i \subseteq \{\text{states} \geq k - i + 1 \text{ steps from error}\}$

Four invariants:

- $F_0 = I$
- $F_i \Rightarrow F_{i+1}$
- $F_i \wedge T \Rightarrow F'_{i+1}$
- Except $F_{k+1}, F_i \Rightarrow P$

\therefore if ever $F_i = F_{i+1}$, F_i is inductive & P is invariant



Essence of IC3

- Continual refinement of over-approximating stepwise sets
 - Until one is inductive
 - Monolithic use of induction
- Generation of clauses as response to backward reachable states
 - Inductive generalization: $c \subseteq \neg s$
(c is inductive relative to a stepwise set)
 - Incremental use of induction

Two Views of IC3

- Prover: Generates predicates from counterexamples
 - From s : state that can reach error
 - To $c \subseteq \neg s$: inductive relative to F_i
 - c proves that s is unreachable in $\leq i + 1$ steps
- Bug finder: Guided backward search
 - Stepwise sets: proximity estimate to initial state

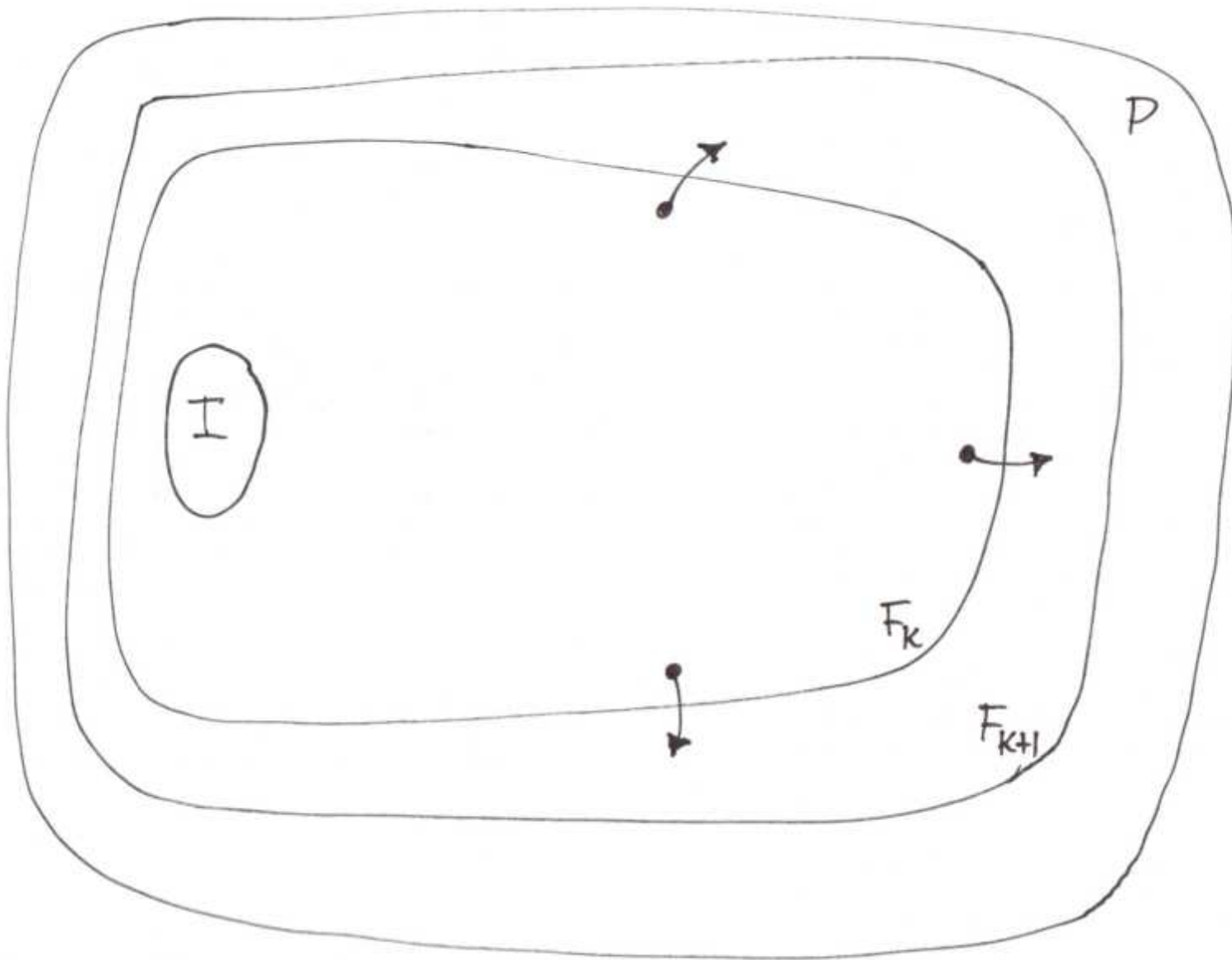
Induction at Top Level

Is P inductive relative to F_k ?

$$F_k \wedge T \Rightarrow P'$$

(Recall: $F_k \Rightarrow P$)

- Possibility #1: Yes
- Conclusion: P is inductive relative to F_k



$$F_k(\wedge P) \wedge T \Rightarrow P'$$

Induction at Top Level

Monolithic behavior (predicate abstraction):

- For i from 1 to k : find largest $C \subseteq F_i$ s.t.

$$F_i \wedge T \Rightarrow C'$$

$$F_{i+1} := F_{i+1} \wedge C$$

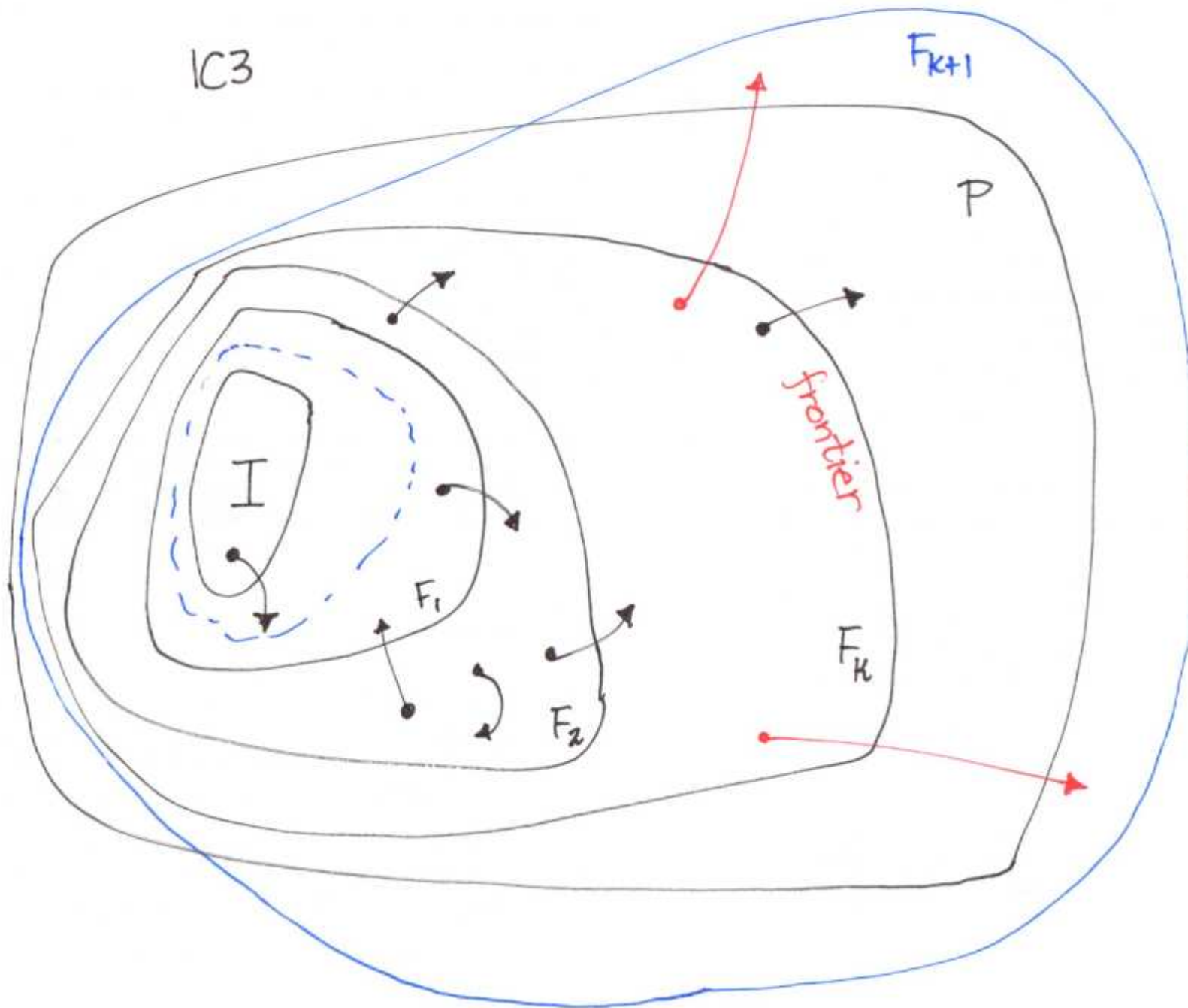
- $F_{k+1} := F_{k+1} \wedge P$
- New frontier: F_{k+1}

If ever $F_i = F_{i+1}$, done: P is invariant.

Counterexample To Induction (CTI)

$$F_k \wedge T \Rightarrow P'$$

- Possibility #2: No
- Conclusion: $\exists F_k$ -state s with error successor
- If s is an initial state, done: P is not invariant
- Otherwise...



Induction at Low Level

Inductive Generalization in IC3

- **Given:** cube s
- **Find:** $c \subseteq \neg s$ such that

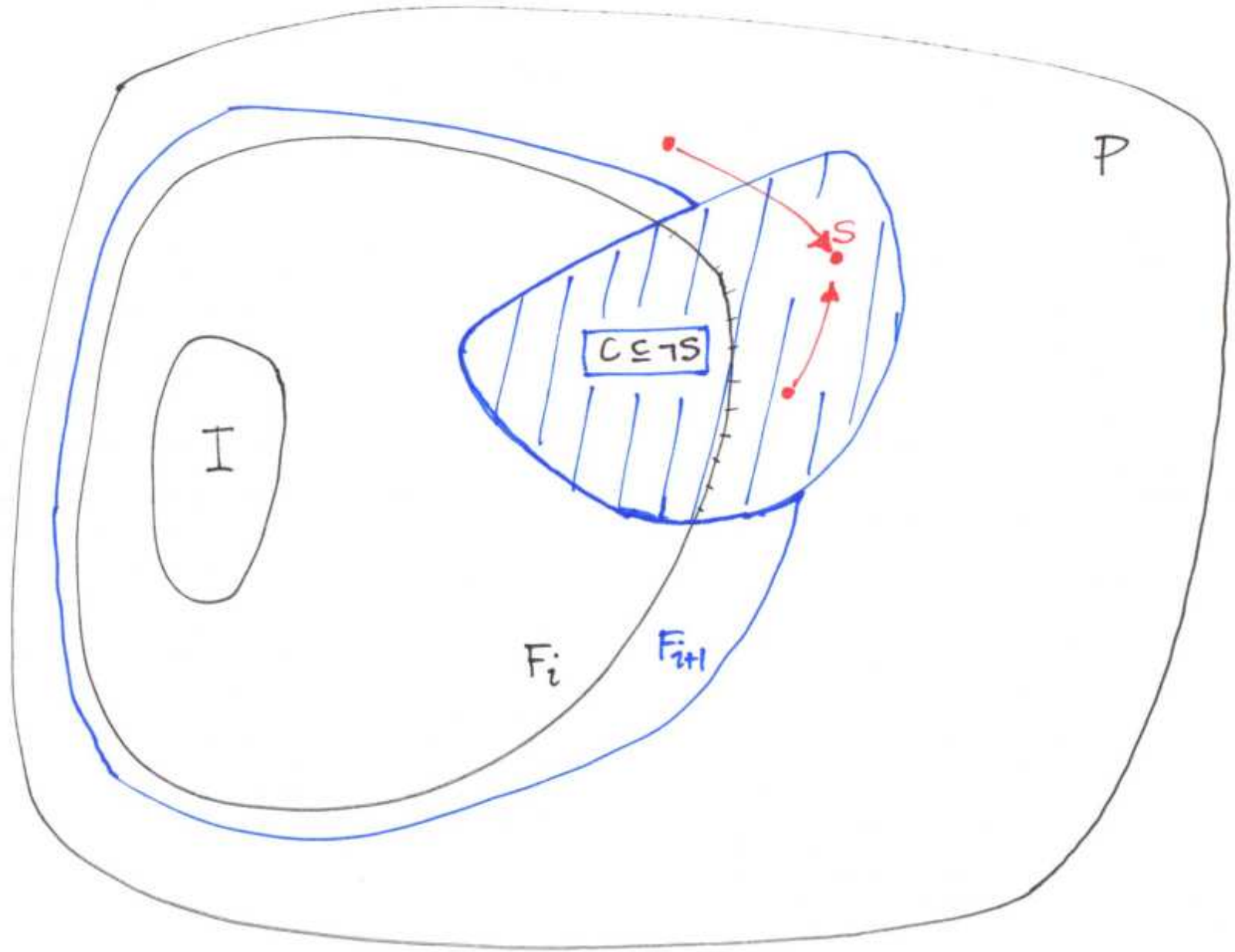
- Initiation:

$$I \Rightarrow c$$

- Consecution (relative to F_i):

$$F_i \wedge c \wedge T \Rightarrow c'$$

- No strict subclause of c is inductive relative to F_i



Inductive Generalization

Addressing CTI s

- Find highest i such that

$$F_i \wedge \neg s \wedge T \Rightarrow \neg s'$$

- Apply inductive generalization:

$$c \subseteq \neg s \quad I \Rightarrow c \quad F_i \wedge c \wedge T \Rightarrow c'$$

- $\therefore F_{i+1} := F_{i+1} \wedge c$ (also update $F_j, j \leq i$)
- If $i < k$, new **proof obligation**:

$$(s, i + 1)$$

“Inductively generalize s relative to F_{i+1} ”

Addressing Proof Obligation (t, j)

SAT query:

$$F_j \wedge \neg t \wedge T \Rightarrow \neg t'$$

If UNSAT:

- Inductive generalization must succeed:

$$c \subseteq \neg t \quad I \Rightarrow c \quad F_j \wedge c \wedge T \Rightarrow c'$$

- $F_{j+1} := F_{j+1} \wedge c$
- Updated proof obligation (if $j < k$): $(t, j + 1)$

Addressing Proof Obligation (t, j)

SAT query:

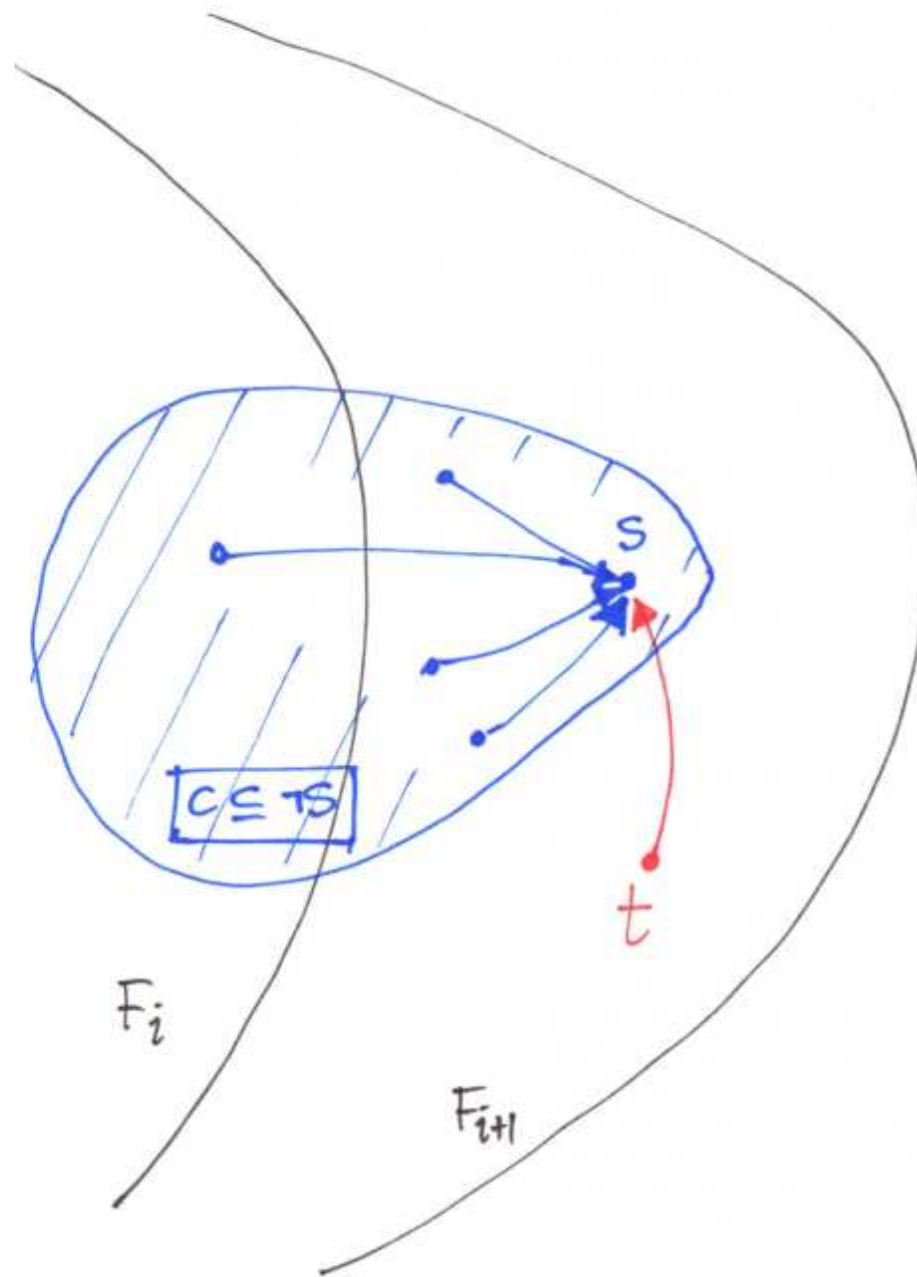
$$F_j \wedge \neg t \wedge T \Rightarrow \neg t'$$

If SAT: New CTI u , treat as before

One of IC3's Insights

Identification of relevant predecessors:

- Why did inductive generalization of s
 - succeed relative to F_i but
 - fail relative to F_{i+1} ?
- Because of some F_{i+1} -state s -predecessor t .
- Analysis at F_i focuses IC3's choice of predecessors at F_{i+1} .



Relevant Predecessor

IC3: A Prover

- Based on CTIs (s), IC3 generates F_i -relative inductive clauses ($c \subseteq \neg s$) to refine F_i 's.
- IC3 propagates clauses to prepare new frontier.
 - Some clauses may be too specific.
 - Their loss can break mutual support.
- As the frontier advances, IC3 considers ever more general situations.
- It eventually finds the real reasons (as truly inductive clauses) that P is invariant.

IC3: A Bug Finder

Suppose:

- $u \rightarrow t \rightarrow s \rightarrow \text{Error}$
- Proof obligations:

$$\{(s, k - 1), (t, k - 2), (u, k - 1)\}$$

That is,

- s must be inductively generalize relative to F_{k-1}
- t must be inductively generalize relative to F_{k-2}
- u must be inductively generalize relative to F_{k-1}

Which proof obligation should IC3 address next?

Guided Search

Two observations:

- u is the “deepest” of the states

$$u \rightarrow t \rightarrow s \rightarrow \text{Error}$$

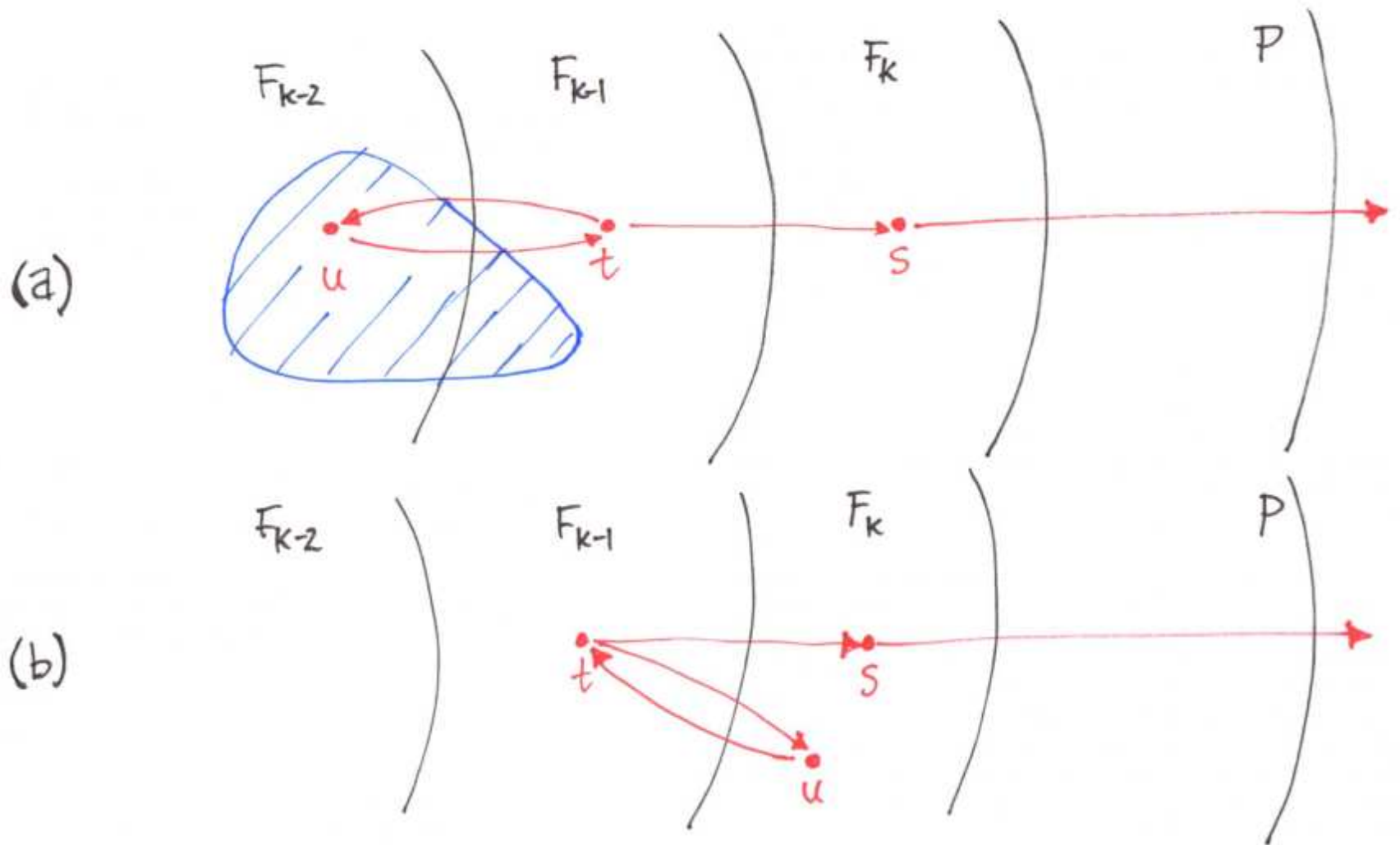
- t is the state that IC3 considers as **likeliest to be closest to an initial state.**

$$\{(s, k - 1), (t, k - 2), (u, k - 1)\}$$

“Proximity metric”

Conclusion: Pursue $(t, k - 2)$ next.

(It also happens to be the correct choice [Bradley 2011].)



$$\{(s, k-1), (t, k-2), (u, k-1)\}$$

Proof Obligations: Guided Search

Incremental, Inductive Verification

IIV Algorithm:

- Constructs concrete **hypotheses**
- Generates intermediate lemmas **incrementally**
- Applies **induction** many times
- **Generalizes** from hypotheses to strong lemmas

After IC3: Refinements

- New heuristic: ternary simulation cube reduction
[Een et al., FMCAD'11]
- Industrial setting: incremental verification
[Chockler et al., FMCAD'11]

Oh, yeah, and a name change: **PDR**
(Thanks, Niklas!)

PDR: Temporal Logics

- FAIR [Bradley et al., FMCAD'11]
 - For ω -regular properties, e.g., LTL
 - Insight: SCC-closed regions can be characterized inductively
- IICTL [Hassan et al., CAV'12]
 - For CTL properties
 - Insight: EX (SAT), EU (IC3), EG (FAIR)
 - Standard traversal of CTL property's parse tree
 - Over- and under-approximating sets
 - Task state-driven refinement

PDR: Infinite-state Systems

- **SMT-based Induction Methods for Timed Systems** [Kindermann et al., arXiv'12]
- **Generalized Property Directed Reachability** [Hoder et al., SAT'12]
 - Boolean push-down systems
 - Linear real arithmetic
- **Software Model Checking via IC3** [Cimatti et al., CAV'12]
 - Explicit handling of CFG
 - Applies IC3 techniques to McMillan's "Lazy Abstraction with Interpolants" [McMillan, CAV'06]

PDR: Handling Proof Obligations

Some presentations use LIFO ordering:

- Trivial correctness; easier to understand
- [Hoder et al., SAT'12], [Cimatti et al., CAV'12]
- Downside: not quite as good?
 - PSPACE-complete (finite-state), so...
 - But: fixed-length counterexamples for K
 - And: not aggressive about mutual induction

Challenges for SAT/SMT

1. Emphasizes incremental calls (100s-1000s/sec)
(FAIR/IICTL: even pushes/pops sets of clauses)
2. Understand effect of solver's choices on PDR
3. Variable ordering:
 - Vital to practical performance
 - Direct core assumptions and lifting
4. IIV: proofs pushed to block (e.g., FAIR)
 - Solver should report whether proof is “useful”
5. Multi-threaded access over core constraints

Conclusions

- Attempted to explain why IC3 works:
 - As a compromise between the **incremental** and **monolithic** strategies
 - Characteristics of previous SAT-based MC
 - As a prover
 - As a bug finder
- Subsequent work: temporal logic, SMT
- Challenges for SAT/SMT